

# Numerical Methods for Computational Science and Engineering

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Recap: Theorem 6.1.6: for every  $f \in C^0([0,1])$   
sequence of Bernstein approximants  $p_n$ , s.t.  
$$\|p_n - f\|_\infty \rightarrow 0 \quad \text{for } n \rightarrow \infty$$

But: slow convergence

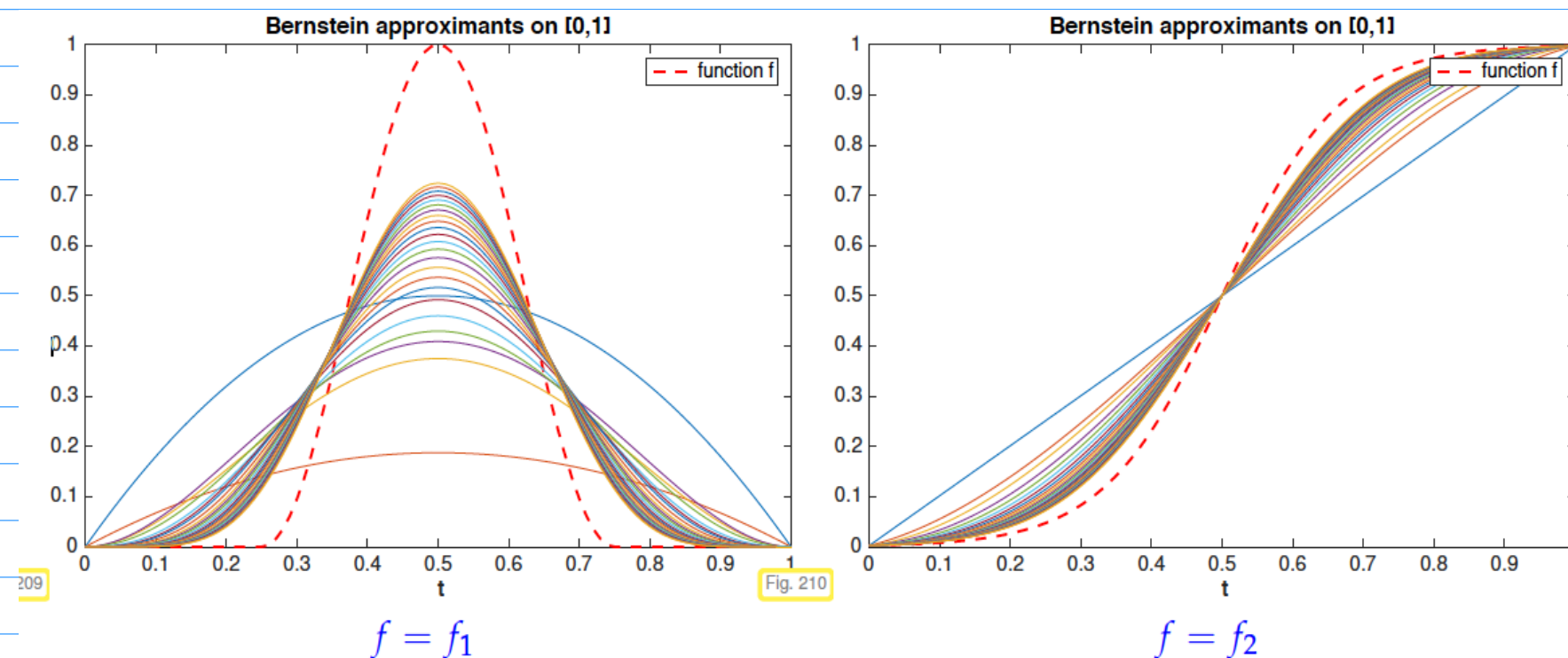
See proof in lecture notes:

$$\forall t \in [0,1]: |f(t) - p_n(t)| \leq (\|f\|_\infty + 1) \varepsilon$$

$$\varepsilon \lesssim \frac{1}{n^\alpha} \quad \alpha < 1$$

Examples:

$$f_1(t) := \begin{cases} 0 & , \text{ if } |2t - 1| > \frac{1}{2} \\ \frac{1}{2}(1 + \cos(2\pi(2t - 1))) & \text{ else,} \end{cases} \quad , \quad f_2(t) := \frac{1}{1 + e^{-12(x-1/2)}}.$$



$$p_n \quad n = 1/2, \dots, 25$$

Are there polynomials that converge faster?

Requires notion of best approximation error

### Definition 6.1.14. (Size of) best approximation error

Let  $\|\cdot\|$  be a (semi-)norm on a space  $X$  of functions  $I \rightarrow \mathbb{K}$ ,  $I \subset \mathbb{R}$  an interval. The (size of the) **best approximation error** of  $f \in X$  in the space  $\mathcal{P}_k$  of polynomials of degree  $\leq k$  with respect to  $\|\cdot\|$  is

$$\text{dist}_{\|\cdot\|}(f, \mathcal{P}_k) := \inf_{p \in \mathcal{P}_k} \|f - p\|.$$

Best possible  $L^\infty$  approximation:

### Theorem 6.1.15. $L^\infty$ polynomial best approximation estimate [Jackson's theorem]

If  $f \in C^r([-1,1])$  ( $r$  times continuously differentiable),  $r \in \mathbb{N}$ , then, for any polynomial degree  $n \geq r$ ,

$$\inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([-1,1])} \leq (1 + \pi^2/2)^r \frac{(n-r)!}{n!} \|f^{(r)}\|_{L^\infty([-1,1])}.$$

$$\inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([-1,1])}$$

- norm of best appr. error
- always exists ( $\mathcal{P}_n$  is finite-dim)
- uniform approximation

Note: This estimate depends on smoothness of  $f$ !

$$r=1: \quad \inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([-1,1])} \leq C \cdot \frac{1}{n} \cdot \|f'\|_{L^\infty([-1,1])}$$

$$r=2: \quad \inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([-1,1])} \leq C^2 \frac{1}{n(n-1)} \|f''\|_{L^\infty([-1,1])}$$

Using Sterling's formula:

$$\sqrt{2\pi} n^{n+1/2} e^{-n} \leq n! \leq e^{1-n} n^{n+1/2}$$

$$\frac{(n-r)!}{n!} \leq \frac{e^{1-(n-r)} (n-r)^{n-r+1/2}}{\sqrt{2\pi} n^{n+1/2} e^{-n}}$$

$$\leq e^{1+r} \frac{1}{\sqrt{2\pi}} \frac{(n-r)^{n-r+1/2}}{n^{n+1/2-r} \cdot n^r}$$

$$\frac{n-r}{n} \leq 1 \rightarrow \leq \tilde{C}(r) \cdot n^{-r}$$

$n-r+1/2 \geq 0$

$$\inf_{p \in P_n} \|f - p\|_{L^\infty([-1,1])} \leq C(r) n^{-r} \|f^{(r)}\|_{L^\infty([-1,1])}$$

for some constant  $C(r)$  depending on  $r$  but independent of  $f$  and  $n$ .

In asymptotic form:

$$\inf_{p \in P_n} \|f - p\|_{L^\infty([-1,1])} = \Theta(n^{-r}) \text{ as } n \rightarrow \infty$$

algebraic convergence in  $n$   
 ↑  
 governed by smoothness of  $f$ !

What if  $f \in C^\infty$ ?

$\forall r \in \mathbb{N}: \exists C_1 = C_1(r)$  s.t.

$$\inf \|f - p\|_\infty \leq C_2(r) n^{-r}$$

$$C_2(r) = C(r) \|f^{(r)}\|_\infty = \left(1 + \frac{\pi^2}{2}\right) \tilde{C}(r) \|f^{(r)}\|_\infty$$

Note:  $\|f^{(r)}\|_\infty$  might not be uniformly bounded in  $r$

Also: Possibly  $\tilde{C}(r) \rightarrow \infty$  as  $r \rightarrow \infty$

(for example:  $r = \frac{n}{2}$ )

$$\tilde{C}(r) = e^{1+r/2} \left(\frac{1}{2}\right)^{\frac{n+1}{2}} \rightarrow \infty \text{ as } n \rightarrow \infty$$

Polynomial approximations on arbitrary intervals  $[a,b]$ :

Approximation schemes on  $[0,1]$  or  $[-1,1]$

can be transformed to appr. scheme on  $[a,b]$ ?

Idea: use affine linear mapping

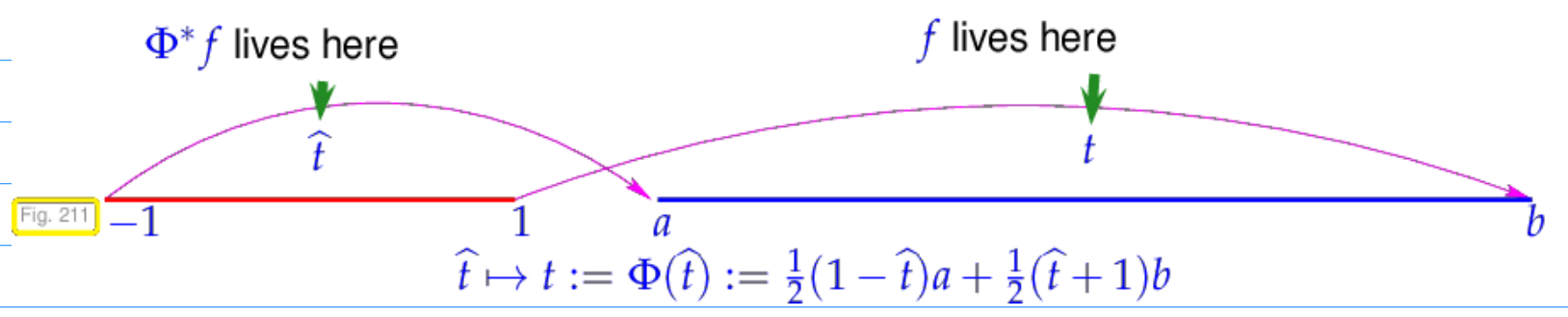
$$\Phi: [-1,1] \rightarrow [a,b]$$

$$\Phi(\hat{t}) = a + \frac{1}{2}(\hat{t}+1)(b-a) \quad \hat{t} \in [-1,1]$$

to introduce affine linear pullback of  $f$ :

$$\Phi^* : C^0([a, b]) \rightarrow C^0([-1, 1]) \quad , \quad \Phi^*(f)(\hat{t}) := f(\Phi(\hat{t})) \quad , \quad -1 \leq \hat{t} \leq 1$$

(6.1.20)



$$\mathbb{F}^* f = f \circ \mathbb{F} \in C^0([-1, 1])$$

**Lemma 6.1.21. Affine pullbacks preserve polynomials**

If  $\Phi^* : C^0([a, b]) \rightarrow C^0([-1, 1])$  is an affine pullback according to (6.1.19) and (6.1.20), then  $\Phi^* : \mathcal{P}_n \rightarrow \mathcal{P}_n$  is a **bijective** linear mapping for any  $n \in \mathbb{N}_0$ .

$\Rightarrow$  For  $\hat{A}$  an appr. scheme on  $C^0([-1, 1])$   
we can define an appr. scheme  $A$  on  $C^0([a, b])$ :

$$A : C^0([a, b]) \rightarrow \mathcal{P}_n$$

$$\hat{A} : C^0([-1, 1]) \rightarrow \mathcal{P}_n$$

$$A := (\mathbb{F}^*)^{-1} \circ \hat{A} \circ \mathbb{F} \quad \mathbb{F}^* : C^0([a, b]) \rightarrow C^0([-1, 1])$$

What is  $(\mathbb{F}^*)^{-1}$  ?

$$f = (\mathbb{F}^*)^{-1} \underbrace{\mathbb{F}^* f}_{f \circ \mathbb{F}} \quad g = \mathbb{F}^* (\mathbb{F}^*)^{-1} g$$

$$(\mathbb{F}^*)^{-1} g = g \circ \mathbb{F}^{-1} \quad f = f \circ \mathbb{F} \circ \mathbb{F}^{-1}$$

$$g = g \circ \mathbb{F}^{-1} \circ \mathbb{F}$$

$$\mathbb{F}^{-1} : [a, b] \rightarrow [-1, 1]$$

$$t \mapsto \frac{2}{b-a} t - \frac{a+b}{b-a}$$

ad Lemma 6.1.21

pullback is linear  $\rightarrow$  suffices to check its action on monomials:

$$\Phi^* \{ t \mapsto t^n \} = \{ \hat{t} \mapsto (a + \frac{1}{2}(\hat{t}+1)(b-a))^n \} \in \mathcal{P}_n$$

$$\Phi^* : \mathcal{P}_n \rightarrow \mathcal{P}_n$$

and restriction of  $(\Phi^*)^{-1}$  to  $\mathcal{P}_n$  has the same property:

$$(\Phi^*)^{-1} : \mathcal{P}_n \rightarrow \mathcal{P}_n$$

Transformation of norms:

#### Lemma 6.1.24. Transformation of norms under affine pullbacks

For every  $f \in C^0([a,b])$  we have

$$\|f\|_{L^\infty([a,b])} = \|\Phi^* f\|_{L^\infty([-1,1])}, \quad \|f\|_{L^2([a,b])} = \sqrt{\frac{b-a}{2}} \|\Phi^* f\|_{L^2([-1,1])}. \quad (6.1.25)$$

clear because  $\Phi^*$  doesn't change point values

$$\begin{aligned} \|f\|_{L^2([a,b])}^2 &= \int_a^b |f(t)|^2 dt = \int_{-1}^1 |f(\Phi(\hat{t}))|^2 \Phi'(\hat{t}) d\hat{t} \\ &= \frac{b-a}{2} \int_{-1}^1 |\Phi^* f(\hat{t})|^2 d\hat{t} \\ &= \frac{b-a}{2} \|\Phi^* f\|_{L^2([-1,1])}^2 \end{aligned}$$

If  $A$  is approximation scheme for  $f \in C^0([a,b])$ :

$$\|f - Af\|_{L^\infty([a,b])} = \|\Phi^* f - \hat{A}(\Phi^* f)\|_{L^\infty([-1,1])}$$

$$\|f - Af\|_{L^2([a,b])} = \sqrt{\frac{b-a}{2}} \|\Phi^* f - \hat{A}(\Phi^* f)\|_{L^2([-1,1])}$$

What can we say about  $\|f^{(r)}\|_{L^\infty([a,b])}$ ?

Employ chain rule:

$$(\Phi^* f)'(\hat{t}) = f'(\Phi(\hat{t})) \cdot \underbrace{\Phi'(\hat{t})}_{\frac{1}{2}(b-a)}$$

$$\Rightarrow (\Phi^* f)^{(r)}(\hat{t}) = \left(\frac{b-a}{2}\right)^r \Phi^* (f^{(r)})(\hat{t})$$

$$\Rightarrow \|\Phi^* f^{(r)}\|_{L^\infty([-1,1])} = \left(\frac{b-a}{2}\right)^r \|f^{(r)}\|_{L^\infty([a,b])} \quad (*)$$

Best polynomial approximation on  $C^0([a,b])$ ?

[Jackson's Thm for arbitrary intervals]

$$\inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([a,b])} = \inf_{p \in \mathcal{P}_n} \|\Phi^* f - \Phi^* p\|_{L^\infty([-1,1])}$$

$$\stackrel{\Phi^*(\mathcal{P}_n) = \mathcal{P}_n}{=} \inf_{p \in \mathcal{P}_n} \|\Phi^* f - p\|_{L^\infty([-1,1])}$$

$$\leq \left(1 + \frac{\pi^2}{2}\right)^r \frac{(n-r)!}{n!} \|(\Phi^* f)^{(r)}\|_{L^\infty([-1,1])}$$

$$\stackrel{(*)}{\leq} \left(1 + \frac{\pi^2}{2}\right)^r \frac{(n-r)!}{n!} \left(\frac{b-a}{2}\right)^r \|f^{(r)}\|_{L^\infty([a,b])}$$

$$\Rightarrow \inf_{p \in \mathcal{P}_n} \|f - p\|_{L^\infty([a,b])} \leq C(r) \left(\frac{b-a}{2n}\right)^r \|f^{(r)}\|_{L^\infty([a,b])}$$

Stirling  
algebraic convergence, but constant also depends on  $|I| = |b-a|$ .



## 6.1.2. Error estimates for polynomial interpolation

### Definition 6.1.32. Lagrangian (interpolation polynomial) approximation scheme

Given an interval  $I \subset \mathbb{R}$ ,  $n \in \mathbb{N}$ , a node set  $\mathcal{T} = \{t_0, \dots, t_n\} \subset I$ , the Lagrangian (interpolation polynomial) approximation scheme  $L_{\mathcal{T}} : C^0(I) \rightarrow \mathcal{P}_n$  is defined by

$$L_{\mathcal{T}}(f) := l_{\mathcal{T}}(\mathbf{y}) \in \mathcal{P}_n \quad \text{with} \quad \mathbf{y} := (f(t_0), \dots, f(t_n))^T \in \mathbb{K}^{n+1}.$$

Behavior as number of nodes is increased:

Different families  $\mathcal{T}_n = \{t_0^{(n)}, \dots, t_n^{(n)}\}$  of nodes

$$\mathcal{T}_n := \left\{ t_j^{(n)} := a + (b-a) \frac{j}{n}, j=0, \dots, n \right\} \subset I$$

For family of polynomial approximation schemes  $\{A_n\}_{n \in \mathbb{N}}$ :

$$\|f - A_n f\| \stackrel{?}{\leq} T(n) \quad \text{for } n \rightarrow \infty$$

here: Lagrange interp. with equidistant nodes

i.e.  $\|f - L_{\mathcal{T}_n} f\| \stackrel{?}{\leq} T(n)$  [bound on interp. error]

Example:  $f(t) = \sin(t)$   $I = [0, \pi]$

Lagrange interp. with equidistant nodes

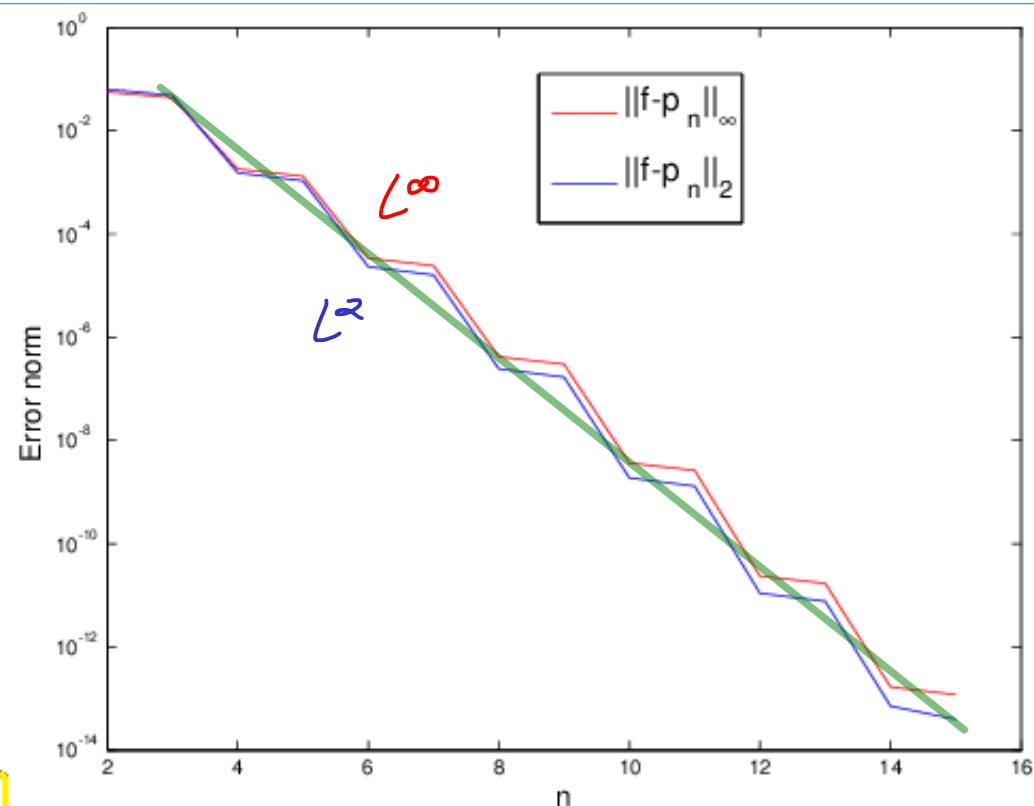


Fig. 212

$\epsilon_n := \|f - L_{\mathcal{T}_n} f\|$  observation:  $\log \epsilon_n \approx C - kn$   $k > 0$   
roughly linear

$$\varepsilon_n \approx e^C e^{-kn}$$

i.e.  $\varepsilon = \Theta(q^n)$  for some  $q \in (0, 1)$ ,  $n \rightarrow \infty$

↑  
exponential convergence

Algebraic convergence:  $\|f - L_{\mathcal{T}}f\| = O(n^{-p})$   
Exponential convergence:  $\|f - L_{\mathcal{T}}f\| = O(q^n)$  for  $n \rightarrow \infty$  ("asymptotic!")

In this example:  $\sin(t) \in C^\infty$

simple Lagrange interpolation is doing much better

than predicted by Jackson's theorem [which is always for fixed smoothness  $r$ ]

Example:  $f(t) = \frac{1}{1+t^2}$   $t \in \mathbb{R}$   $I = [-5, 5]$   
again equidistant Lagrange interpolation

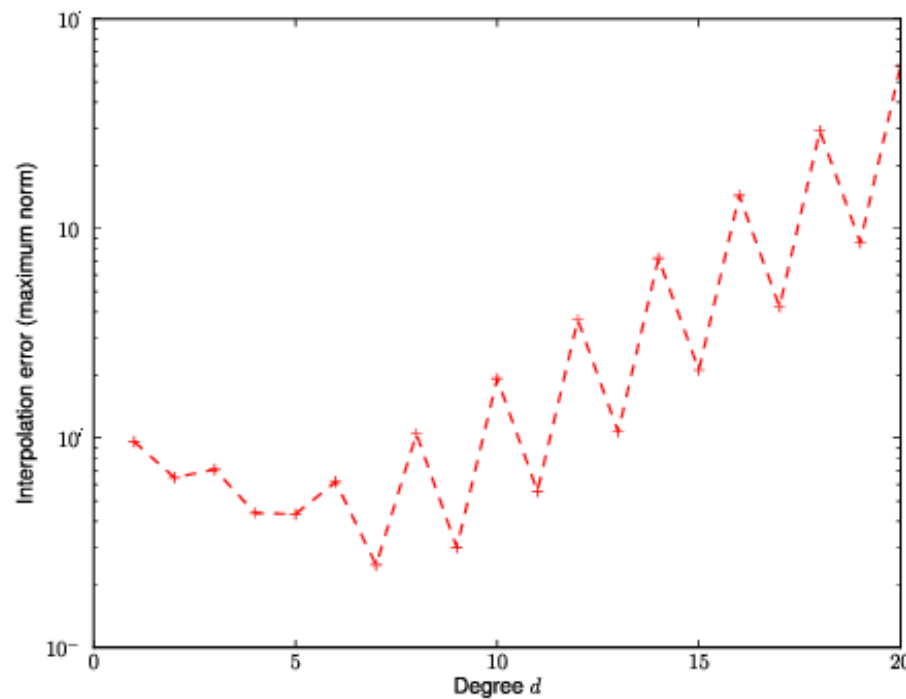


Fig. 214

Approximate  $\|f - L_{\mathcal{T}_n f}\|_\infty$  on  $[-5, 5]$

Recall:

Runge's phenomenon



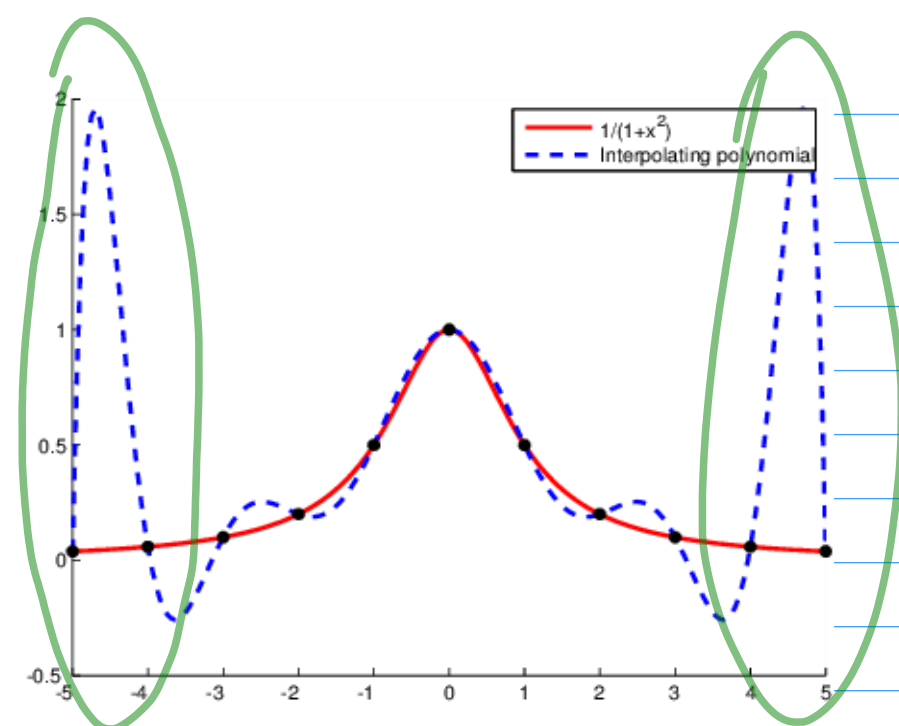


Fig. 213

Interpolating polynomial,  $n = 10$

Can we understand the interpolation error?

**Theorem 6.1.44. Representation of interpolation error** [?, Thm. 8.22], [?, Thm. 37.4]

We consider  $f \in C^{n+1}(I)$  and the Lagrangian interpolation approximation scheme ( $\rightarrow$  Def. 6.1.32) for a node set  $\mathcal{T} := \{t_0, \dots, t_n\} \subset I$ . Then, for every  $t \in I$  there exists a  $\tau_t \in ]\min\{t, t_0, \dots, t_n\}, \max\{t, t_0, \dots, t_n\}[$  such that

$$f(t) - L_{\mathcal{T}}(f)(t) = \underbrace{\frac{f^{(n+1)}(\tau_t)}{(n+1)!}}_{=: c} \cdot \underbrace{\prod_{j=0}^n (t - t_j)}_{=: \omega(t)} \quad (6.1.45)$$

"nodal polynomial"

Proof: Fix  $t \in I \setminus \mathcal{T}$   $f \in C^{n+1}(I)$

We can choose  $c \in \mathbb{R}$  s.t.

$$\underbrace{f(t)}_{\text{fixed}} - L_{\mathcal{T}}f(t) - c\omega(t) = 0$$

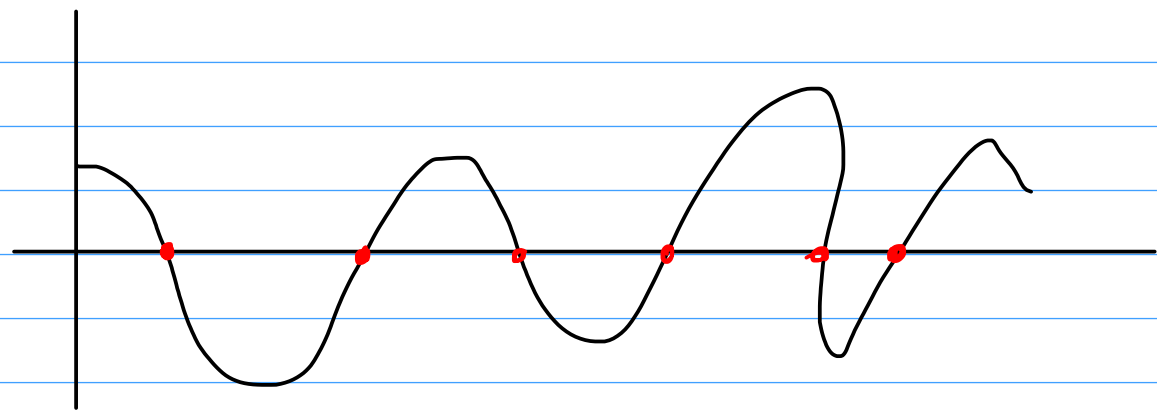
Define function:

$$\varphi(x) := \underbrace{f(x) - L_{\mathcal{T}}f(x)}_{\in \mathcal{P}_n} - \underbrace{c\omega(x)}_{\in \mathcal{P}_{n+1}} \in C^{n+1}(I)$$

How many zeros does  $\varphi$  have at least?

$$\left. \begin{aligned} \varphi(t_j) &= 0 && (I.C) && j=0, \dots, n && \rightarrow n+1 \text{ zeros} \\ \varphi(t) &= 0 && (\text{by def.}) \end{aligned} \right\}$$

$\Rightarrow n+2$  distinct zeros



$\varphi'$  has at least  $n+1$  distinct zeros [mean value thm]

$\varphi''$   $n$  distinct zeros

$\varphi^{(n+1)}(x)$  at least 1 zero name it  $\tau_t$

$$\varphi^{(n+1)}(x) = f^{(n+1)}(x) - 0 - c(n+1)!$$

and

$$\varphi^{(n+1)}(\tau_t) = 0 = f^{(n+1)}(\tau_t) - c(n+1)!$$

$$c = \frac{f^{(n+1)}(\tau_t)}{(n+1)!} \quad \square$$

Use this pointwise estimate to get a global estimate:

$$\text{Thm. 6.1.44} \Rightarrow \|f - L_T f\|_{L^\infty(I)} \leq \frac{\|f^{(n+1)}\|_{L^\infty(I)}}{(n+1)!} \max_{t \in I} |(t-t_0) \cdots (t-t_n)| \quad (6.1.50)$$

Example of  $\sin(t)$  on  $[0, \pi]$ :

$$\|f^{(n)}\|_\infty \leq 1$$

$$\Rightarrow \|f - L_T f\|_{L^\infty(I)} \leq \frac{1}{(n+1)!} \max_{t \in [0, \pi]} |t \cdot (t - \frac{\pi}{n}) \cdot (t - \frac{2\pi}{n}) \cdots (t - \pi)|$$

↑  
equidistant nodes & Lagrange interp. extremal at  $\approx \frac{\pi}{2n}$

$$\leq \frac{1}{(n+1)!} \left| \frac{\pi}{2n} \left( \frac{\pi}{2n} - \frac{\pi}{n} \right) \cdots \left( \frac{\pi}{2n} - \pi \right) \right|$$

$$\leq \frac{1}{(n+1)!} \left( \frac{\pi}{n} \right)^{n+1} \underbrace{\left| \frac{1}{2} \left( \frac{1}{2} - 1 \right) \left( \frac{1}{2} - 2 \right) \cdots \left( \frac{1}{2} - n \right) \right|}_{\leq n!}$$

$$\leq \frac{1}{n+1} \left(\frac{\pi}{n}\right)^{n+1}$$

↑  
faster than exponential

Example:  $f(t) = \frac{1}{1+t^2}$  on  $I = [-5, 5]$

$$\|f^{(n)}\|_{L^\infty([-5, 5])} \sim 2^n n! \quad \text{for } n \rightarrow \infty$$

Estimate (6.1.50) no longer guarantees convergence  
(i.e. blow-up possible)

[ RHS of (6.1.50) roughly:

$$\begin{aligned} &\sim 2^{n+1} (n+1)! \frac{1}{(n+1)!} n! \left(\frac{5}{n}\right)^{n+1} 2^n \\ &= n! \left(\frac{20}{n}\right)^n \cdot \frac{10}{n} \end{aligned}$$

By Stirling this grows exponentially:

$$\underbrace{n^{n+1/2} \sqrt{2\pi} e^{-n} \left(\frac{20}{n}\right)^n \frac{10}{n}}_{= \left(\frac{20}{e}\right)^n \sqrt{2\pi} \frac{10}{n^{1/2}}} \leq n! \left(\frac{20}{n}\right)^n \frac{10}{n}$$

$\left(\frac{20}{e}\right)^n$  exp. growth!

Note: there is also an  $L^2$ -estimate for  $f \in C^{n+1}(I)$ .

and  $\mathcal{T} := \{t_0, \dots, t_n\} \in I$

$$\|f - L_{\mathcal{T}} f\|_{L^2(I)} \leq \frac{2^{(n-1)/4} |I|^{n+1}}{\sqrt{n! (n+1)!}} \|f^{(n+1)}\|_{L^2(I)}$$

### 6.1.3 Chebyshev interpolation

Recall: RHS of (6.1.50) depended on

$\|f^{(n+1)}\|_{L^\infty(I)}$  and

$\max_{t \in I} |w(t)|$

← here we have some choice

we can't control

New task: Given  $n$  and  $I$

find nodes  $t_0, \dots, t_n$  s.t.

$\|w\|_{L^\infty(I)}$  minimal!

Recall:  $w(t) = \prod_{j=0}^n (t - t_j) \in \mathcal{P}_{n+1}$   
leading coeff. is 1

Equivalent problem:

Find  $q \in \mathcal{P}_{n+1}$  with leading coeff. 1  
s.t.  $\|q\|_{L^\infty(I)}$  minimal

This implies that  $q$  has  $n+1$  zeros in  $I$ :

Take  $I = [-1, 1]$  and  $q(t) := \prod_{j=0}^n (t - t_j)$

with  $t_0 < -1$ ,

Then, define

$$p(t) := (t+1) \cdot (t-t_1) \dots (t-t_n)$$

$$\Rightarrow |p(t)| < |q(t)| \quad \forall t \in I = [-1, 1]$$

$$\Rightarrow \|p\|_{L^\infty(I)} < \|q\|_{L^\infty(I)} \quad \hookrightarrow \text{to } q \text{ having min norm}$$

Recipe once  $q$  is found:

Take nodes  $t_0, \dots, t_n$  to be the zeros of  $q$ .

As we will see: For  $q \in \mathcal{P}_{n+1}$  with  $\|q\|_{L^\infty(I)}$  min.:

$\|q\|_{L^\infty}$  will be attained at edges of the boundary & local extrema of  $q$  (in abs. value):

Can we find the minimizing  $q$ ?

→ Chebyshev polynomials

Definition 6.1.76. Chebyshev polynomials → [?, Ch. 32]

The  $n^{\text{th}}$  Chebyshev polynomial is  $T_n(t) := \cos(n \arccos t)$ ,  $-1 \leq t \leq 1, n \in \mathbb{N}$ .

$$|T_n(t)| \leq 1 \quad \forall t \in [-1, 1]$$

Are these actual polynomials?

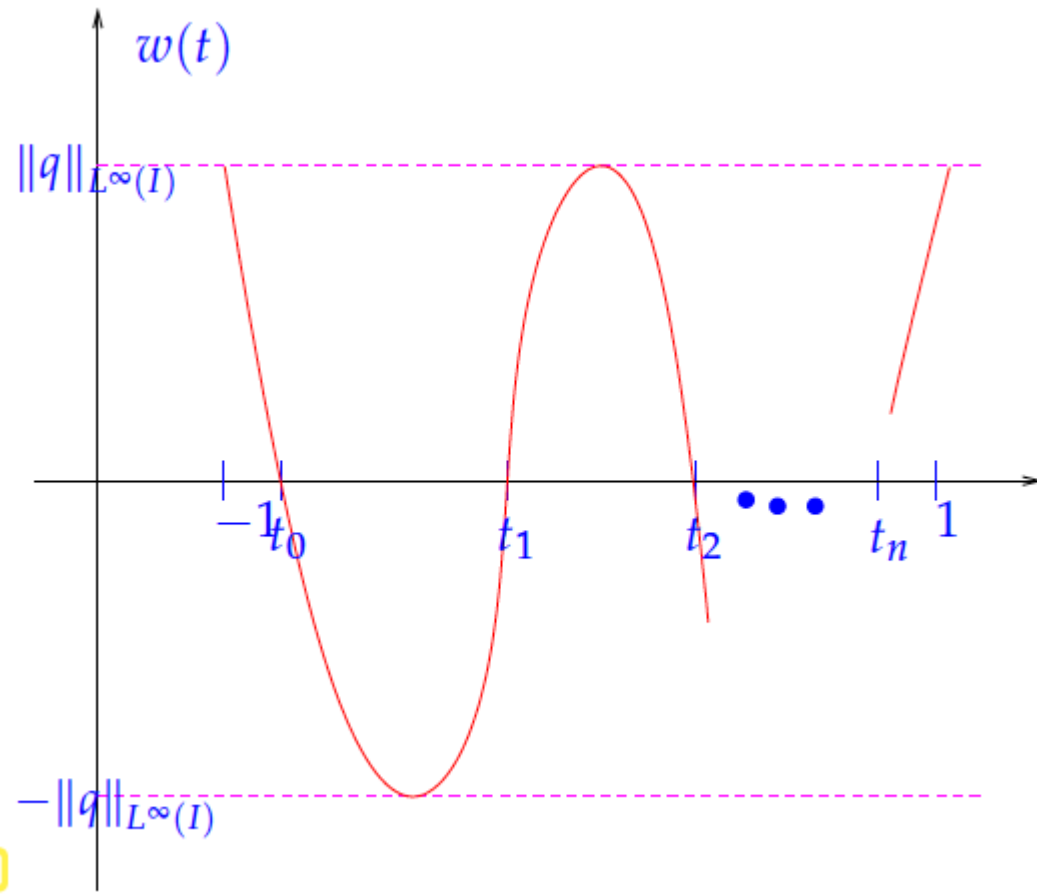
Theorem 6.1.77. 3-term recursion for Chebyshev polynomials → [?, (32.2)]

The function  $T_n$  defined in Def. 6.1.76 satisfy the 3-term recursion

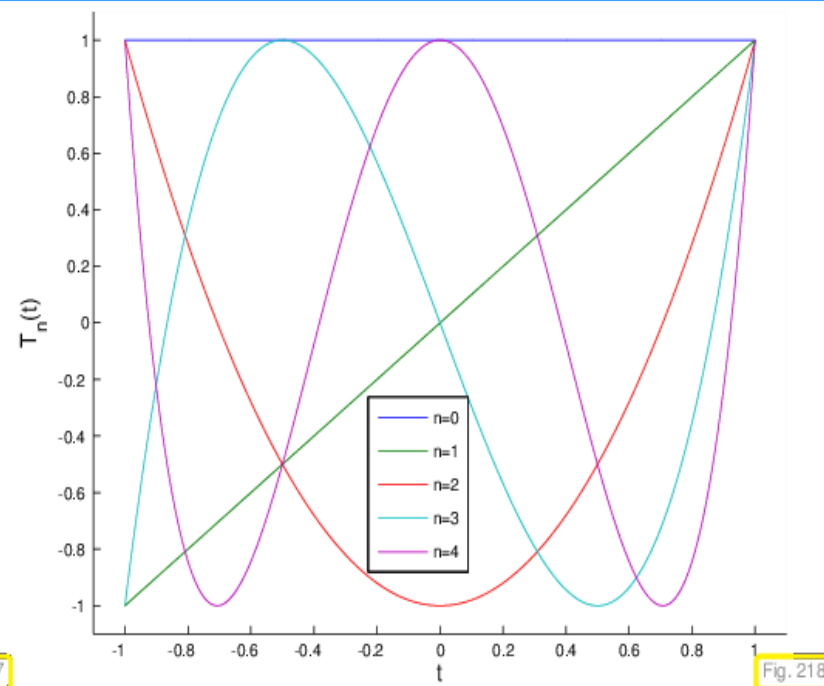
$$T_{n+1}(t) = 2tT_n(t) - T_{n-1}(t), \quad T_0 \equiv 1, \quad T_1(t) = t, \quad n \in \mathbb{N}. \quad (6.1.78)$$

If  $T_{n-1} \in \mathcal{P}_{n-1}, T_n \in \mathcal{P}_n \Rightarrow T_{n+1} \in \mathcal{P}_{n+1}$

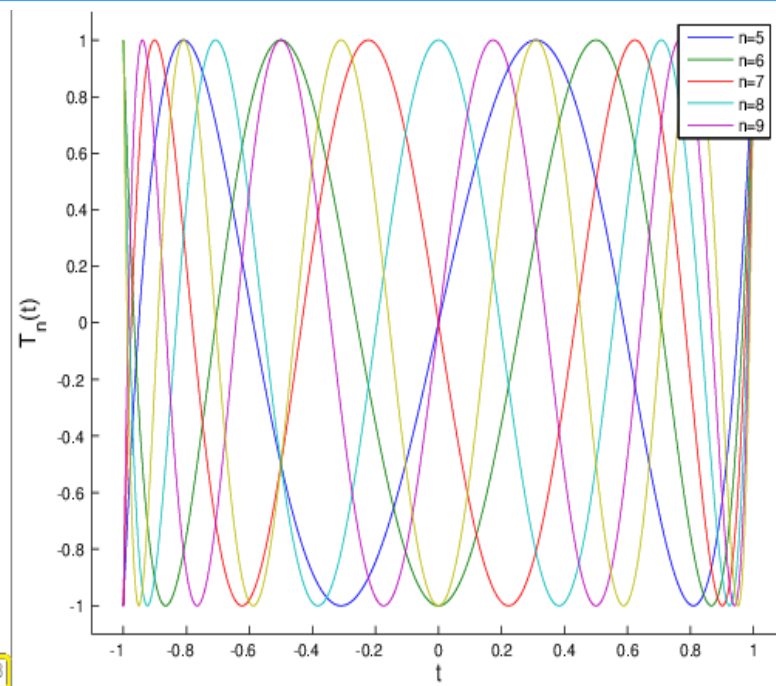
$T_0 \in \mathcal{P}_0, T_1 \in \mathcal{P}_1 \Rightarrow$  by induction:  $T_n \in \mathcal{P}_n \quad \forall n \in \mathbb{N}$ .



Note: Goal here is an a-priori choice of nodes without using information on  $f$ !  
A-posteriori methods also exist.



Chebyshev polynomials  $T_0, \dots, T_4$



Chebyshev polynomials  $T_5, \dots, T_9$

$$\arccos(t_j) = \frac{2j+1}{2n} \pi$$

$$t_j = \cos\left(\frac{2j+1}{2n} \pi\right) \quad \text{Chebyshev nodes}$$

Is this choice optimal?

**Theorem 6.1.82. Minimax property of the Chebyshev polynomials** [?, Section 7.1.4.], [?, Thm. 32.2]

The polynomials  $T_n$  from Def. 6.1.76 minimize the supremum norm in the following sense:

$$\|T_n\|_{L^\infty([-1,1])} = \inf\{\|p\|_{L^\infty([-1,1])} : p \in \mathcal{P}_n, p(t) = 2^{n-1}t^n + \dots\}, \quad \forall n \in \mathbb{N}.$$

$$T_2(t) = 2t^2 - 1$$

$$T_3(t) = 4t^3 - 2t$$

From the recursion: leading coefficient of  $T_n$ :  $2^{n-1}$

Zeros of  $T_n(t)$ :

$$n \arccos(t_j) = \frac{2j+1}{2} \pi \quad j=0, \dots, n$$

Proof:

① local extrema of  $T_n$ ?

$$|\cos x| = 1 \iff x = j\pi \quad \text{for some } j \in \mathbb{Z}$$

$$|T_n(\tilde{t}_j)| = 1 \iff n \arccos(\tilde{t}_j) = j\pi$$

$$\tilde{t}_j = \cos\left(\frac{j\pi}{n}\right) \quad j=0, \dots, n$$



$\Rightarrow n+1$  local extrema of  $T_n$  in  $I = [-1, 1]$ .

Suppose  $\exists q \in \mathcal{P}_n$  with leading coeff.  $2^{n-1}$  s.t.

$$\|q\|_{L^\infty(I)} < \|T_n\|_{L^\infty(I)} = 1$$

$\Rightarrow$  • If  $\tilde{t}_j$  local minimum of  $T_n$ :  $T_n(\tilde{t}_j) = -1$

$$\Rightarrow (T_n - q)(\tilde{t}_j) < 0$$

• If  $\tilde{t}_j$  local maximum of  $T_n$ :  $T_n(\tilde{t}_j) = 1$

$$\Rightarrow (T_n - q)(\tilde{t}_j) > 0$$

# of local extrema  $n+1$ :

$T_n - q$  changes sign at least  $n$  times  
on  $[-1, 1]$

$T_n - q$  has at least  $n$  zeros

$T_n$  and  $q$  have the same leading coefficient

$$\Rightarrow T_n - q \in \mathcal{P}_{n-1}$$

$$T_n - q \equiv 0 \quad \hookrightarrow \text{to } \|q\|_{L^\infty(I)} < \|T_n\|_{L^\infty(I)} \quad \square$$

This result implies: nodal polynomial

$$\omega(t) = \prod_{j=0}^n (t - t_j)$$

with  $\|\omega\|_{L^\infty(I)}$  minimal is found by

$$\omega(t) := 2^{-n} T_{n+1}(t)$$

and  $t_j$  are zeros of  $T_{n+1}$  (Chebyshev nodes)

$$t_j = \cos\left(\frac{2j+1}{2n} \pi\right) \quad j=0, \dots, n$$

With choice  $w(t) = 2^{-n} T_{n+1}(t)$   
 what can we say about interpolation error?

Now  $\|w\|_{L^\infty([-1,1])} = 2^{-n}$

$f \in C^{n+1}([-1,1])$

$\|f - L_n f\|_{L^\infty([-1,1])} \leq \frac{2^{-n}}{(n+1)!} \|f^{(n+1)}\|_{L^\infty([-1,1])}$

On arbitrary intervals  $[a,b]$ :

transform nodes with affine linear transformation:

$[-1,1] \rightarrow [a,b]$   
 $\hat{t} \mapsto a + \frac{b-a}{2} (\hat{t} + 1) \quad \hat{t} \in [-1,1]$

Distribution of Chebyshev nodes: more dense at boundaries as  $n$  increases

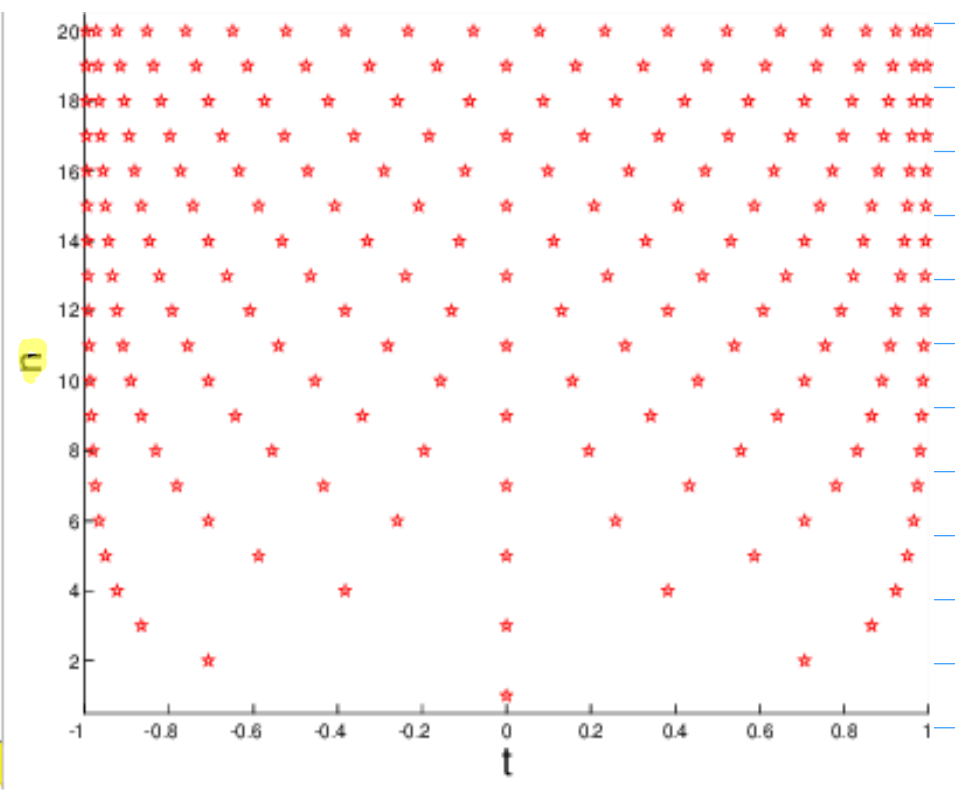


Fig. 219

$I = [-1, 1]$

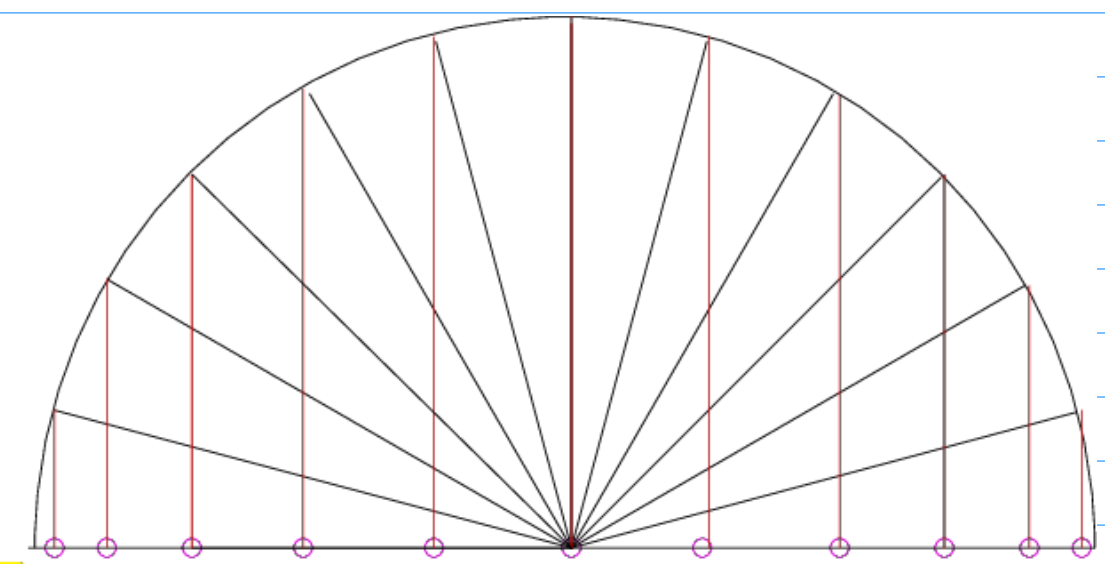
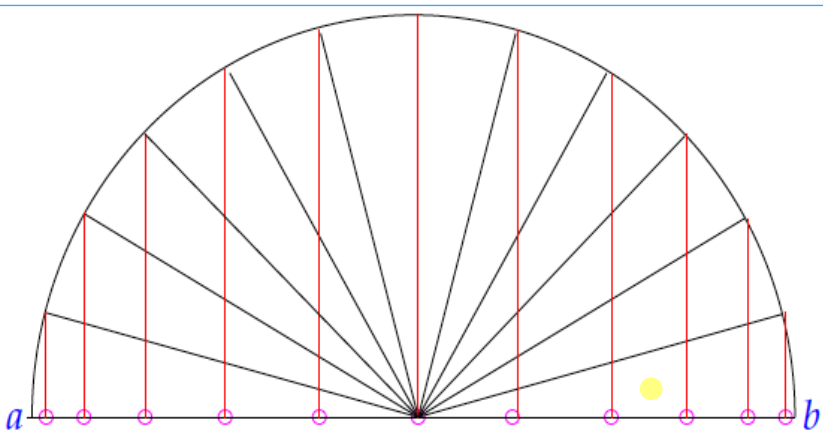


Fig. 220

equidistant on circle  
 $\rightarrow$  more dense at edges



The **Chebyshev nodes** in the interval  $I = [a, b]$  are

$$t_k := a + \frac{1}{2}(b-a) \left( \cos\left(\frac{2k+1}{2(n+1)}\pi\right) + 1 \right), \quad (6.1.87)$$

$k = 0, \dots, n.$

Interpolation error estimate for general  $I = [a, b]$ :

$$\|f - I_T(f)\|_{L^\infty(I)} = \overset{\text{pullback}}{\| \hat{f} - I_{\hat{T}}(\hat{f}) \|_{L^\infty([-1,1])}} \leq \frac{2^{-n}}{(n+1)!} \left\| \frac{d^{n+1}\hat{f}}{d\hat{t}^{n+1}} \right\|_{L^\infty([-1,1])}$$

$$\leq \frac{2^{-2n-1}}{(n+1)!} |I|^{n+1} \|f^{(n+1)}\|_{L^\infty(I)}. \quad (6.1.88)$$

much better estimate than for equidistant nodes

But: still possible divergence if  $\|f^{(n)}\|_{L^\infty(I)}$  grows too fast.

Example:  $\frac{1}{1+t^2} \quad t \in [-5, 5]$

$$\|f^{(n+1)}\|_{L^\infty([-5,5])} \sim 2^{n+1} (n+1)!$$

still RHS blows up:  $\frac{2^{-2n-1}}{(n+1)!} 10^{n+1} 2^{n+1} (n+1)!$

$$= 5^n \cdot 10$$

↑  
exp. growth

BUT: not necessarily divergence

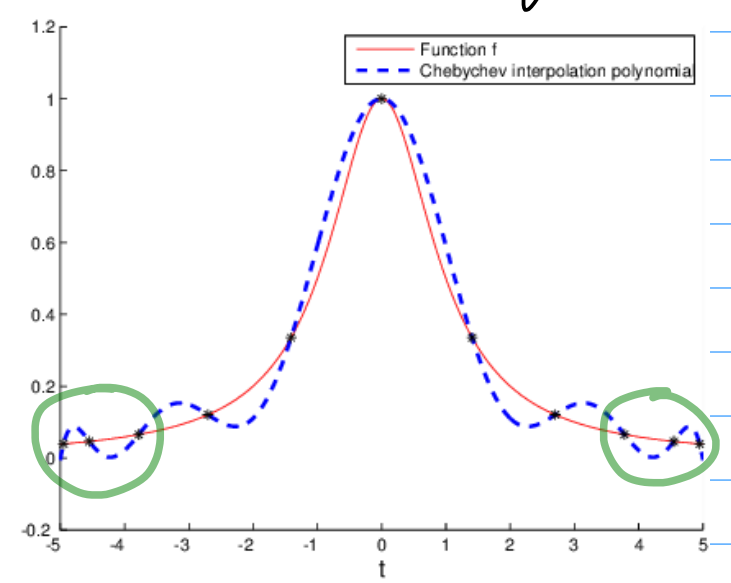


Fig. 222

Chebyshev nodes  $n=10$

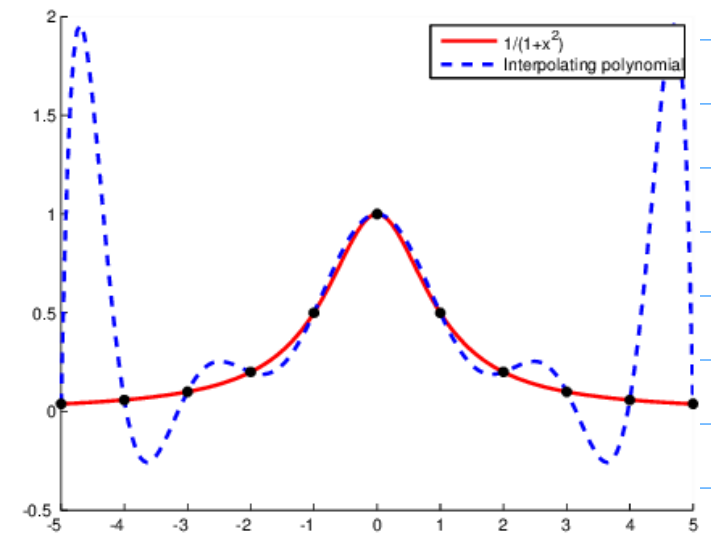


Fig. 221

Equidistant nodes

Estimates independent of  $\|f^{(n)}\|_\infty$

Concept: Lebesgue constant [quality measure for polyn. interp. scheme]

Given  $\mathcal{I} = \{t_0, \dots, t_n\}$

$$I_{\mathcal{I}}: \mathbb{R}^{n+1} \rightarrow \mathcal{P}_n$$

$$I_{\mathcal{I}}((f(t_0), \dots, f(t_n))) = L_{\mathcal{I}}f$$

Lebesgue constant:  $\lambda_{\mathcal{I}} := \sup_{\gamma \in \mathbb{R}^{n+1}} \frac{\|I_{\mathcal{I}}\gamma\|_\infty}{\|\gamma\|_\infty}$

$$L_{\mathcal{I}}: C^0(I) \rightarrow \mathcal{P}_n$$

$$f \mapsto L_{\mathcal{I}}f = I_{\mathcal{I}}(\underbrace{(f(t_0), \dots, f(t_n))}_{=z})$$

$$\|L_{\mathcal{I}}f\|_{L^\infty(I)} = \|I_{\mathcal{I}}(z)\|_{L^\infty(I)}$$

$$= \frac{\|I_{\mathcal{I}}(z)\|_\infty}{\|z\|_\infty} \cdot \|z\|_\infty \leq \lambda_{\mathcal{I}} \|z\|_\infty$$

$$\leq \lambda_{\mathcal{I}} \|f\|_{L^\infty(I)}$$

$$\Rightarrow \|L_{\mathcal{I}}f\|_{L^\infty(I)} \leq \lambda_{\mathcal{I}} \|f\|_{L^\infty(I)}$$

[ $\lambda_{\mathcal{I}}$ : norm of the operator  $I_{\mathcal{I}}$  in  $L^\infty$ -sense]

More explicitly:  $\lambda_{\mathcal{I}} = \max_{t \in I} \sum_{j=0}^n |L_j(t)|$

Lagrange polynomials

Implication for interpolation error:

$$\|f - L_{\mathcal{I}}f\|_{L^\infty(I)} = \|f - p - (L_{\mathcal{I}}f - p)\|_{L^\infty(I)} \quad \text{for any } p \in \mathcal{P}_n$$

$$\stackrel{L_{\mathcal{I}}p = p}{=} \|f - p - \underbrace{L_{\mathcal{I}}(f-p)}_{\text{linear}}\|_{L^\infty(I)}$$

$$\begin{aligned} &\leq \|f-p\|_\infty + \underbrace{\|L_\gamma(f-p)\|_\infty}_{\substack{\uparrow \\ \text{triangle} \\ \text{ineq.}}} \\ &\leq \|f-p\|_\infty + \lambda_\gamma \|f-p\|_\infty \\ &= (1+\lambda_\gamma) \|f-p\|_\infty \end{aligned}$$

$$\|f - L_T f\|_{L^\infty(I)} \leq (1 + \lambda_T) \inf_{p \in P_n} \|f - p\|_{L^\infty(I)} \quad \forall f \in C^0(I), \quad (6.1.61)$$

best approximation error

interpolation error is at most  $1 + \lambda_\gamma$  worse than the best approximation error!

(6.1.61) is a special case of Lebesgue's lemma in approximation theory:

given  $(X, \|\cdot\|)$  normed vector space,  $U$  subspace of  $X$  and  $P$  is a linear projection on  $U$ ,

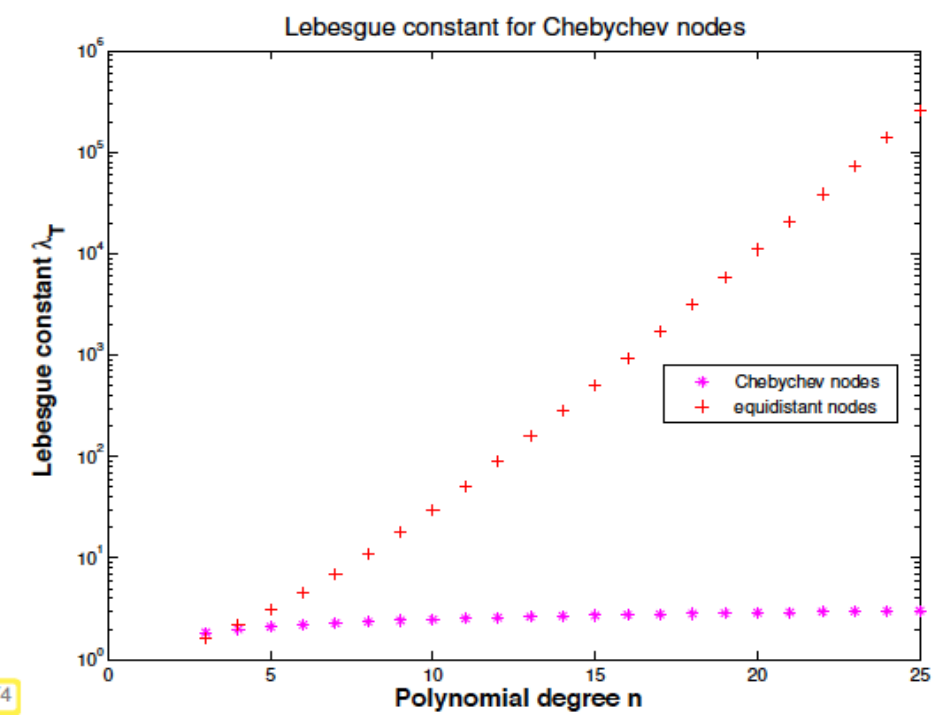
$$P: X \rightarrow U$$

$$\text{Then } \|x - P x\| \leq (1 + \|P\|) \inf_{u \in U} \|x - u\| \quad \forall x \in X$$

Quantify  $\lambda_\gamma$  ?

Equidistant nodes:  $\lambda_\gamma \geq C e^{n/2}$  (at least exp. growth)

Chebyshev nodes:  $\lambda_\gamma \leq \frac{2}{\pi} \log(n+1) + 1$  (at most log growth)



← equidistant nodes

lin-log-plot

← Chebyshev nodes

Combine general estimate (6.1.61) with Jackson's thm:

$$\|f - L_{\sigma} f\|_{L^{\infty}([-1,1])} \leq (1 + \lambda_{\sigma}) \inf_{p \in \mathcal{P}_n} \|f - p\|_{L^{\infty}([-1,1])}$$

$$\leq (1 + \lambda_{\sigma}) \left(1 + \frac{\pi^2}{2}\right)^r \frac{(n-r)!}{n!} \|f^{(r)}\|_{L^{\infty}([-1,1])}$$

for  $f \in C^r([-1,1])$

For Chebyshev interpolation:

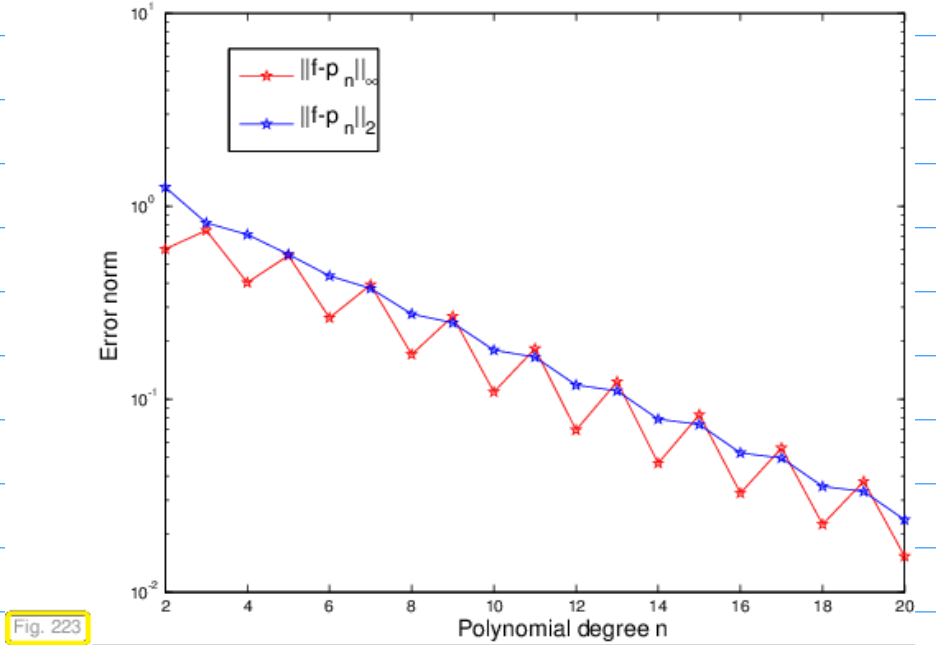
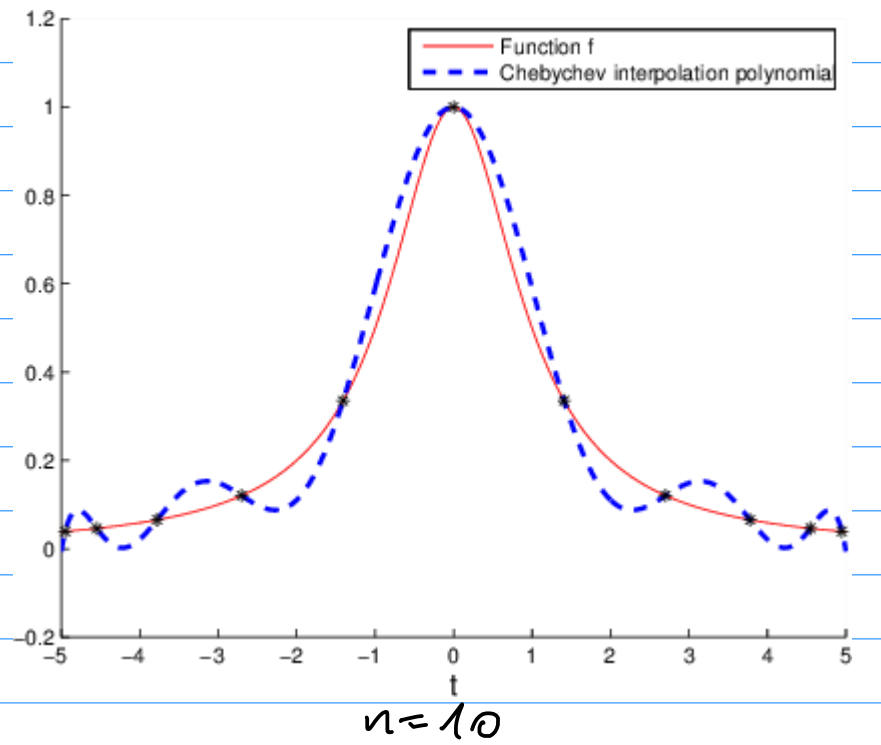
$$\|f - L_{\tau} f\|_{L^{\infty}([-1,1])} \leq \underbrace{(2/\pi \log(1+n) + 2)}_{1 + \lambda_{\sigma}} \underbrace{\left(1 + \frac{\pi^2}{2}\right)^r \frac{(n-r)!}{n!}}_{\leq C(r) n^{-r}} \|f^{(r)}\|_{L^{\infty}([-1,1])}$$

alg. convergence

for fixed  $r$  and  $n \rightarrow \infty$ .

Examples with Chebyshev interpolation:

①  $f(t) = \frac{1}{1+t^2} \quad t \in [-5, 5]$



roughly: exp. convergence  
mostly true when  $f \in C^{\infty}$



② Hat function only  $C^0$

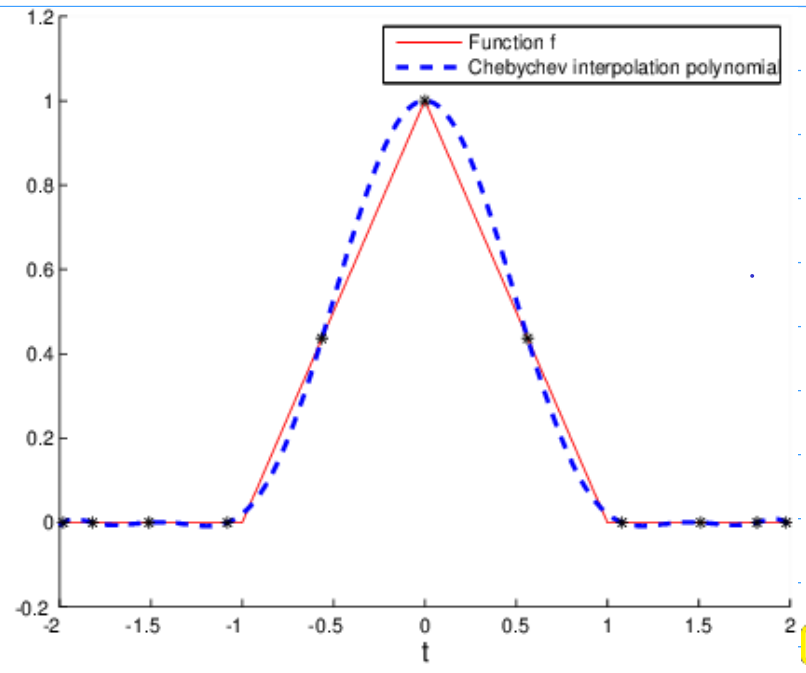


Fig. 224

$n=10$

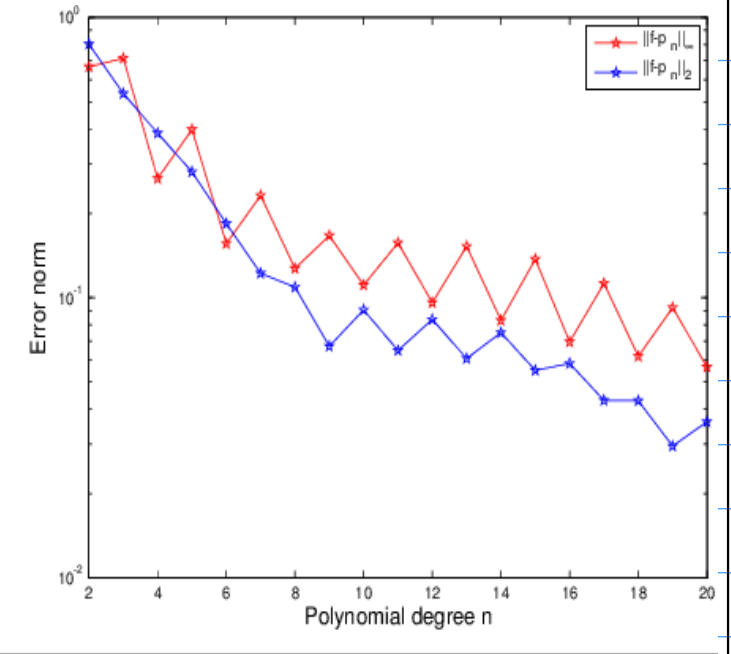


Fig. 225

still: convergence (not exp.)

③  $f(t) = \begin{cases} \frac{1}{2}(1 + \cos(\pi t)) & |t| < 1 \\ 0 & |t| \in [1, 2] \end{cases}$   $I = [-2, 2]$

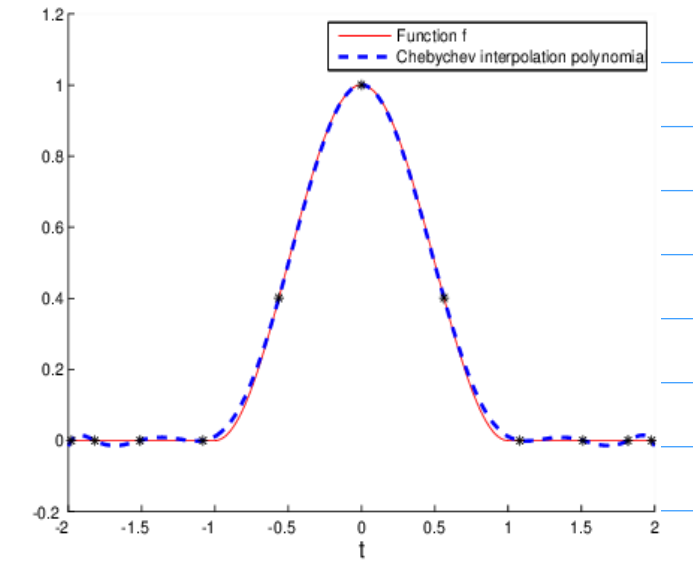


Fig. 227

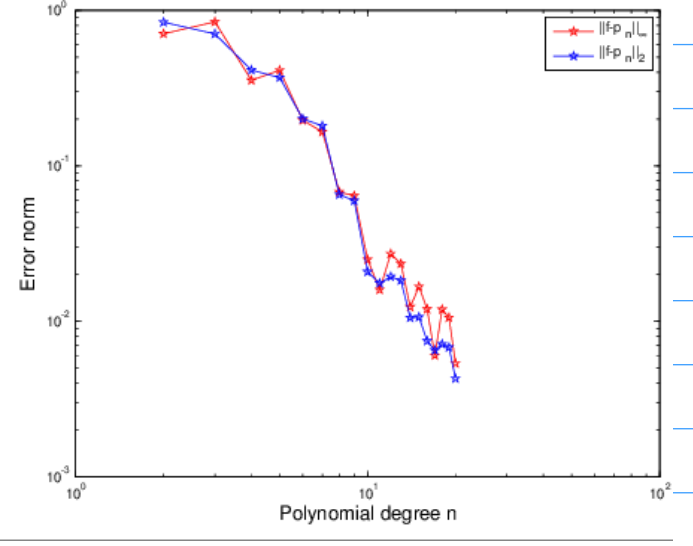


Fig. 228

~ exp. convergence

### 6.1.3.3 Implementation of Chebyshev interpolation

Idea: interpolant  $p \in \mathcal{P}_n$  as

$$p(t) = \sum_{j=0}^n \alpha_j T_j(t)$$

$\uparrow$   
 Chebyshev polynomial  
 basis of  $\mathcal{P}_n$

$\deg T_j = j \Rightarrow \{T_0, \dots, T_n\}$  basis of  $\mathcal{P}_n$

① Efficient evaluation of  $p$  given coefficients  $\alpha_j$

Idea: Use the 3-term recursion of Chebyshev polynomials:

$$T_j(t) = 2t T_{j-1}(t) - T_{j-2}(t) \quad j=2, \dots, n$$

$$p(x) = \sum_{j=0}^{n-1} \alpha_j T_j(x) + \alpha_n T_n(x)$$

$$\stackrel{(6.1.78)}{=} \underbrace{\sum_{j=0}^{n-1} \alpha_j T_j(x)}_{*} + \alpha_n (2x T_{n-1}(x) - T_{n-2}(x))$$

$$= \sum_{j=0}^{n-3} \alpha_j T_j(x) + (\alpha_{n-2} - \alpha_n) T_{n-2}(x) + (\alpha_{n-1} + 2x\alpha_n) T_{n-1}(x)$$

$$(*) : \sum_{j=0}^{n-3} \alpha_j T_j(x) + \alpha_{n-2} T_{n-2}(x) + \alpha_{n-1} T_{n-1}(x)$$

→ another Chebyshev expansion

$$p(x) = \sum_{j=0}^{n-1} \tilde{\alpha}_j T_j(x) \quad \text{with} \quad \tilde{\alpha}_j = \begin{cases} \alpha_j + 2x\alpha_{j+1} & , \text{ if } j = \underline{n-1}, \\ \alpha_j - \alpha_{j+2} & , \text{ if } j = \underline{n-2}, \\ \alpha_j & \text{ else.} \end{cases} \quad (6.1.103)$$

**C++11 code 6.1.104: Recursive evaluation of Chebychev expansion (6.1.101)**

```
2 // Recursive evaluation of a polynomial  $p = \sum_{j=1}^{n+1} a_j T_{j-1}$  at point  $x$ 
3 // based on (6.1.103)
4 // IN : Vector of coefficients  $a$ 
5 // evaluation point  $x$ 
6 // OUT: Value at point  $x$ 
7 double reclenshaw(const VectorXd& a, const double x) {
8     const VectorXd::Index n = a.size() - 1;
9     if (n == 0) return a(0); // Constant polynomial
10    else if (n == 1) return (x*a(1) + a(0)); // Value  $\alpha_1 * x + \alpha_0$ 
11    else {
12        VectorXd new_a(n);
13        new_a << a.head(n - 2), a(n - 2) - a(n), a(n - 1) + 2*x*a(n);
14        return reclenshaw(new_a, x); // recursion
15    }
16 }
```

Complexity:  $\mathcal{O}(n)$ .