

**Question 1:**

Let  $X$  be a topological space and  $f: [0, 1] \rightarrow X$  be a path (recall, this means  $f$  is continuous). Show that  $f + f^{-1}$  is a 1-boundary in two different ways:

1. combining lemmas proven in class,
2. finding an explicit 2-chain whose boundary is  $f + f^{-1}$ .

**Solution:**

**Solution 1:** Observe that  $f + f^{-1}$  is a boundary if and only if  $\llbracket f \rrbracket + \llbracket f^{-1} \rrbracket = 0$  in  $H_1(X)$ . By Lemma 2 of the lecture we have  $\llbracket f \rrbracket + \llbracket f^{-1} \rrbracket = \llbracket f * f^{-1} \rrbracket$ . Let  $x = f(0)$ . Then  $f * f^{-1}$  is homotopic relative to  $x$  to the constant path  $c_x$ . By Lemma 1 we have  $\llbracket f * f^{-1} \rrbracket = \llbracket c_x \rrbracket$ . By Lemma 0,  $\llbracket c_x \rrbracket = 0$ , which concludes the first proof.

**Solution 2:** Let  $\sigma: \Delta_2 \rightarrow X$  be the singular simplex defined as follows:

$$\sigma(\lambda_0 e_0 + \lambda_1 e_1 + \lambda_2 e_2) = f(\lambda_1),$$

and let  $\sigma_0: \Delta_2 \rightarrow X$  be the simplex with constant value  $f(0)$ . Let  $\rho = \sigma + \sigma_0$ . We claim that  $\partial\rho = f + f^{-1}$ , which provides the result. Firstly, observe that  $\partial\sigma = f + f^{-1} - c_{f(0)}$ . Indeed, let  $s_0, s_1$  be a basis for  $\Delta_1$ . Then

$$\sigma^{(0)}: ((1-t)s_0 + ts_1) \mapsto (0e_0 + (1-t)e_1 + te_2) \mapsto f(1-t)$$

$$\sigma^{(1)}: ((1-t)s_0 + ts_1) \mapsto ((1-t)e_0 + 0e_1 + te_2) \mapsto f(0)$$

$$\sigma^{(2)}: ((1-t)s_0 + ts_1) \mapsto ((1-t)e_0 + te_1 + 0e_2) \mapsto f(t).$$

Reasoning as in Question 2, we obtain that  $\partial\sigma_0 = c_{f(0)}$ . Thus,  $\partial\rho = f^{-1} + f - c_{f(0)} + c_{f(0)} = f^{-1} + f$ . *Attention!* It is not true that  $\partial\sigma = f + f^{-1}$ , because it is not true that  $f^{-1} + f + c_{f(0)} = f^{-1} + f$  in  $\Delta_1(X)$ . The equality only holds in  $H_1(X)$ .

## Question 2:

Let  $X$  be a topological space,  $x \in X$  a point and for each  $n$  consider the constant singular simplex  $c_n: \Delta_n \rightarrow X$  with constant value  $x$  (that is the map that sends every point of  $\Delta_n$  to  $x$ ).

When is  $c_n$  a boundary?

### Solution:

First, we compute when  $c_n$  is a cycle. Since  $c_n$  is the constant simplex on  $x$ , all of its faces are  $c_{n-1}$ , i.e. the constant function  $\Delta_{n-1} \rightarrow \{x\}$ . Then for  $n \geq 1$  we have

$$\partial c_n = \sum_{i=0}^n (-1)^i c_n^{(i)} = \sum_{i=0}^n (-1)^i c_{n-1} = \begin{cases} 0 & \text{if } n \text{ is odd;} \\ c_{n-1} & \text{if } n \text{ is even.} \end{cases}$$

Thus, for  $n$  odd we have that  $c_n = \partial c_{n+1}$  showing that it is a boundary. For  $n$  even, we have that  $c_n$  is not even a cycle, in particular not a boundary. Thus  $c_n$  is a boundary precisely when it is a cycle and precisely in odd degrees.

## Question 3:

Let  $F_2$  be the free group on two generators. Show that the abelianization of  $F_2$  is  $\mathbb{Z}^2$ . As a remark, observe that the proof also works to show that the abelianization of  $F_n$  is  $\mathbb{Z}^n$ .

### Solution:

Let  $F_2$  be generated by  $a, b$  and let  $\mathbb{Z}^2$  be generated by  $x, y$ . Recall that  $F_2$  is defined as follows: the elements of  $F_2$  are words in the letters  $a, b, a^{-1}, b^{-1}$ , under the equivalence relation  $uaa^{-1}v \sim ua^{-1}av \sim ubb^{-1}v \sim ub^{-1}bv \sim uv$ , where  $u$  and  $v$  are any words.

Given an element  $g \in F_2$ , let  $\phi_a(g)$  be the sum with sign of all the occurrences of the letter  $a$  in a word  $w$  representing  $g$ . Note that the equivalence relation preserves the exponent sum, so if  $w' \sim w$  both represent the element  $g$ , then their exponent sum will be the same. Let  $\phi_b$  be defined similarly for  $b$ . Consider the map

$$\begin{aligned} \phi: F_2 &\rightarrow \mathbb{Z}^2 \\ g &\mapsto (\phi_a(g), \phi_b(g)). \end{aligned}$$

We claim that it is a group homomorphism. Let  $g, h \in F_2$ , choose a representative  $v$  for  $g$  and  $w$  for  $h$ , and consider  $vw$ . It is clear that the number of occurrences of the letter  $a$  in  $vw$  precisely coincide with the number of occurrences of  $a$  in  $v$  plus the ones in  $w$ , and similarly for  $a^{-1}$ . Since  $vw$  is a representative of  $gh$ , and since the exponent sum does not depend on the choice of representative, we get

$$\phi_a(gh) = \phi_a(g) + \phi_a(h),$$

and similarly for  $\phi_b$ . Thus  $\phi$  is a group homomorphism. We claim that the kernel of  $\phi$  coincides with the commutator subgroup  $[F_2, F_2]$ . Indeed, this would imply  $F_2/[F_2, F_2] \cong \mathbb{Z}^2$ .

The kernel of  $\phi$  consists of those elements of  $F_2$  which are both in the kernel of  $\phi_a$  and  $\phi_b$ , namely those elements  $g$  that can be represented by a word which has the same number of occurrences of  $a, a^{-1}$  and the same number of occurrences of  $b, b^{-1}$ , for example  $aab^{-1}a^{-1}a^{-1}b$ . Let's call this set  $\Sigma$ . Observe that  $[F_2, F_2] \leq \Sigma$ . Indeed,  $[F_2, F_2]$  is generated by elements of the form  $ghg^{-1}h^{-1}$ , which is not hard to see that belong to  $\Sigma$ . Now, let  $g \in F_2$  be such that  $g \in \Sigma$  and let  $w$  be a shortest representative of  $g$ . We will show by induction on the length of  $w$  that it can be written as a product of commutators. First observe that  $w$  needs to have even length.

*Base case:* If the length of  $w$  is less or equal to 4, and  $w$  is minimal in its equivalence class, then  $w$  is either empty or a commutator (this can be done by hand by just listing all words of length at most 4 with zero exponent sum).

*Induction step:* The strategy is going to be the following. We will find  $z$  which represents an element of  $[F_2, F_2]$  such that, after performing reduction,  $wz$  is a shorter word with exponent sum zero. The induction hypothesis then yields  $wz \in [F_2, F_2]$ . Since we picked  $z \in [F_2, F_2]$ , we conclude that  $w \in [F_2, F_2]$ , which will complete the proof.

To simplify notation, let  $A = a^{-1}$  and  $B = b^{-1}$ . Consider  $w$  with zero exponent sum. If all the letters of  $w$  are in  $\{a, A\}$ , the fact that  $w$  has zero exponent sum and it is the shortest representative yield that  $w$  is the empty word, and similarly for  $\{b, B\}$ . Then, up to exchanging the roles of  $a$  and  $A$ , we can write  $w$  as  $\tau av$ , for some words  $t, u$  and  $v$ . Let  $z = [v^{-1}a, u^{-1}]$ . Then: We have

$$wz = \tau av[v^{-1}a, u^{-1}] = \tau avv^{-1}au^{-1}Avu = tvu$$

It is straightforward to verify that  $tvu$  has still exponent sum zero, concluding the proof.