Mathematics of Machine Learning Homework 11

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Try to solve the questions before looking to the answers. Every item must be proved rigorously. Starred problems are harder.

Problem 1

We say that two unit vectors $u, v \in \mathbb{R}^d$ are ε -almost orthogonal if $|\langle u, v \rangle| \le \varepsilon$. A set of vectors in which every vector is ε -almost orthogonal to each other is called ε -almost orthogonal set. For example, for $\varepsilon = 0$, we recover the usual notion of orthogonal vectors and the canonical basis is an 0-almost orthogonal set. Clearly, the maximal size of an 0-orthogonal set in \mathbb{R}^d is *d*. The goal of this exercise is to show that $\varepsilon > 0$, the size of an ε -almost orthogonal set is exponentially large in *d* and discuss its relevance for frame theory. The proof strategy relies on the probabilistic method, a nonconstructive strategy to show the existence of certain object by using probability. The following inequality will be useful (you do not need to prove it), if X_1, \ldots, X_n are i.i.d Rademacher random variables, i.e, taking values 1 and -1 with probability $\frac{1}{2}$ each, then

$$\mathbb{P}(\sum_{i=1}^n X_i \geq t) \leq e^{\frac{-t^2}{2n}}.$$

- (a) Let $\varphi \in \mathbb{R}^d$ be a random vector drawn as follows: Each entry is generated independently at random following a Rademacher distribution. Consider *N* i.i.d copies of the random vector φ , namely $\varphi_1, \ldots, \varphi_N$. Prove the following statement: If $N = \lfloor e^{d\epsilon^2/4} \rfloor$, then the probability that set of normalized vectors $\{\frac{1}{\sqrt{d}}\varphi_1, \ldots, \frac{1}{\sqrt{d}}\varphi_N\}$ is an ε -almost orthogonal set is positive (strictly larger than 0).
- (b) Conclude that letter "a" implies in the existence of ε -almost orthogonal set with size exponentially large in *d*.
- (c) Give an interpretation of this fact in terms of mutual coherence of frames in the sense of Challenge 11.4 in the lecture notes.
- (d) * Improve the estimate of letter "a" from $N = \lfloor e^{d\varepsilon^2/4} \rfloor$ to $N = \lfloor ce^{d\varepsilon^2/2} \rfloor$, where *c* is an absolute constant less than one. (Hint: Compute the expected number of unordered pair that violates the condition of ε -almost orthogonality, then delete one vector from each pair).

Problem 2

The goal of this exercise is to solve the Challenge 12.4 in the lecture notes. Recall that a random graph $G \sim \mathcal{G}(n, p)$ is a graph of *n* vertices generated by placing each possible edge independently at random with probability *p*. We use the standard asymptotic notation, $f(n) \ll g(n)$ if $\lim_{n\to\infty} \frac{f(n)}{g(n)} = 0$, also $g(n) \gg f(n)$ if $\lim_{n\to\infty} \frac{g(n)}{f(n)} = \infty$. Moreover f(n) = o(1) if $\lim_{n\to\infty} f(n) = 0$. Finally we define $p := \frac{\lambda \log n}{n}$ for some constant $\lambda > 0$.

- (a) Prove that if $\lambda \le 1 c$, where c > 0 is an absolute constant, then the graph *G* has an isolated vertex with probability 1 o(1). (Hint: Consider a random variable X_n to be the number of isolated vertices in the graph *G*, you need to prove that the probability of X_n be larger than zero is 1 o(1). Use the inequality $\mathbb{P}(X_n = 0) \le \frac{\operatorname{Var} X_n}{\mathbb{P} X^2}$).
- (b) Now observe the following: A graph is disconnected if and only if thre exits a set of k nodes such that $k \leq \lfloor \frac{n}{2} \rfloor$ and there is no edge connecting the set of k nodes with the complement set of n k nodes. Use this fact to prove that if $\lambda \geq 1 + c$ for an absolute constant c > 0, then the graph is connected with probability 1 o(1). (Hint: Use union bound twice).

Problem 3

Let \mathcal{A} and \mathcal{B} be two family of events with finite VC-dimension. Prove the following facts:

- (a) If $C = \mathcal{A} \cup \mathcal{B}$, then $S_C(n) \leq S_A(n) + S_B(n)$.
- (b) If $C = \mathcal{A} \cap \mathcal{B}$, then $S_C(n) \leq S_{\mathcal{A}}(n)S_{\mathcal{B}}(n)$.