

1. Graphs and automorphisms of trees

Def. 1.1:

- A graph X is a tuple consisting of a set of vertices $X^0 \neq \emptyset$, a set of edges X^1 and three maps $\alpha: X^1 \rightarrow X^0$, $\omega: X^1 \rightarrow X^0$, $\bar{\cdot}: X^1 \rightarrow X^1$ ("beginning", "end" and "inverse" of an edge) such that $\forall e \in X^1$

$$\bar{\bar{e}} = e, \quad \bar{e} \neq e \quad \text{and} \quad \alpha(\bar{e}) = \omega(e).$$
 The vertices $\alpha(e)$ and $\omega(e)$ are called the initial and terminal vertices of the edge e .
- A graph is finite if $X^0 \cup X^1$ is finite.

- A (graph) morphism $p: X \rightarrow Y$ between graphs X and Y is a map that sends vertices to vertices, edges to edges and satisfies

$$p(\alpha(e)) = \alpha(p(e)), \quad p(\omega(e)) = \omega(p(e)), \\ p(\bar{e}) = \overline{p(e)}.$$

It is called an isomorphism if it is bijective (on vertices and edges). An automorphism is an isomorphism from a graph to itself.

For $x \in X^0, y \in Y^0$, we also write $p: (X, x) \rightarrow (Y, y)$ if $p(x) = y$ (and we want to emphasize this).

- The star of a vertex $x \in X^0$ is the set

$$st(x) := \{e \in X^1 \mid \alpha(e) = x\}.$$

The valence of x is

$$val(x) := |st(x)|.$$

- A graph is oriented if from each pair $\{e, \bar{e}\}, e \in X^1,$

• A graph is oriented if from each pair $\{e, \bar{e}\}$, $e \in X$, one element is chosen. This edge is called positively oriented. We write X_+ for the set of positively oriented edges and $X_- = X \setminus X_+$ for its complement, the negatively oriented edges.

• A sequence $l = e_1 e_2 \dots e_n$ of edges of a graph X is called a path of length n if

$$\omega(e_i) = \alpha(e_{i+1}) \quad \forall 1 \leq i \leq n-1.$$

In this case, l is called a path from $\alpha(e_1)$ to $\omega(e_n)$ ("initial" and "terminal" vertex of l).

The path l is closed if $\alpha(e_1) = \omega(e_n)$.

We consider any vertex $v \in X^0$ as a path of length 0 from v to itself.

A path l is called reduced if it has length 0 or if

$$l = e_1 \dots e_n \text{ with } e_{i+1} \neq \bar{e}_i \quad \forall 1 \leq i \leq n-1.$$

• A graph X is connected if for all $v, w \in X^0$, there is a path from v to w .

A circuit in X is a subgraph isomorphic to C_n for some $n \in \mathbb{N}$.

A tree is a connected graph that does not contain a circuit.

Lemma 1.3: If X is a connected graph and T is a maximal subtree of X (with respect to inclusion), then T contains all vertices of X .

Def 1.4: A reduced path in a tree is called a geodesic.

Lemma 1.5: If X is a tree and X_1, X_2 are disjoint subtrees, then there is a unique geodesic with initial vertex in X_1 , terminal vertex in X_2 and all edges outside X_1 and X_2 .

Def. 1.6: Let X be a tree and τ an automorphism of X .

- For $v, w \in X^0$, denote by $[v, w]$ the (unique) geodesic from v to w . Its length is denoted by $d(v, w)$.
- τ acts without inversions if for all $e \in X^1$, we have $\tau(e) \neq \bar{e}$.
- For $x \in X^0 \cup X^1$, we also write $x^\tau := \tau(x)$.

The translation length of τ is defined as

$$|\tau| := \min_{v \in X^0} d(v, v^\tau).$$

If $|\tau| = 0$, we define \tilde{e} to be the subgraph of X consisting of all $x \in X^0 \cup X^1$ such that $x^\tau = x$.

If $|\tau| > 0$, define \tilde{e} as the minimal subtree of X that contains $\{x \in X^0 \mid d(x, x^\tau) = |\tau|\}$.

Theorem 1.7: Let X be a tree and τ an automorphism of X .

i) If $|\tau| = 0$, then \tilde{e} is a tree. Let $v \in X^0$ and let $w \in (\tilde{e})^0$ be a vertex such that $d(v, w)$ is minimal. Then $d(v, w) = d(v^\tau, w)$ and the concatenation of $[v, w]$ and $[w, v^\tau]$ is the geodesic $[v, v^\tau]$ connecting v and v^τ .

ii) If $|\tau| > 0$ and τ acts without inversions, then \tilde{e} is isomorphic to \mathbb{C}^∞ and τ acts on \mathbb{C}^∞ by translation of distance $|\tau|$.

Let $v \in X^0$ and let $w \in \tilde{e}^0$ be a vertex such that $d(v, w)$ is minimal. Then $[v, v^\tau] \cap \tilde{e} = [w, w^\tau]$ and $d(v, v^\tau) = |\tau| + 2 \cdot d(v, w)$.