

Lemma 10.6: Let  $e \in X^1$  and  $k \in \mathbb{N}$ .

There are only finitely many  $k$ -separators that contain  $e$ .  
 Pf.: We prove the statement by induction on  $k$ .

$k=1$  is trivial.

$k \rightarrow k+1$ : Let  $\xrightarrow{e} \xrightarrow{u}$  be an edge of  $X$ .

If there is no  $(k+1)$ -separator that contains  $e$ ,  
 there is nothing to prove.

Otherwise,  $X - \{e\}$  is connected because  $k+1 \geq 2$ .



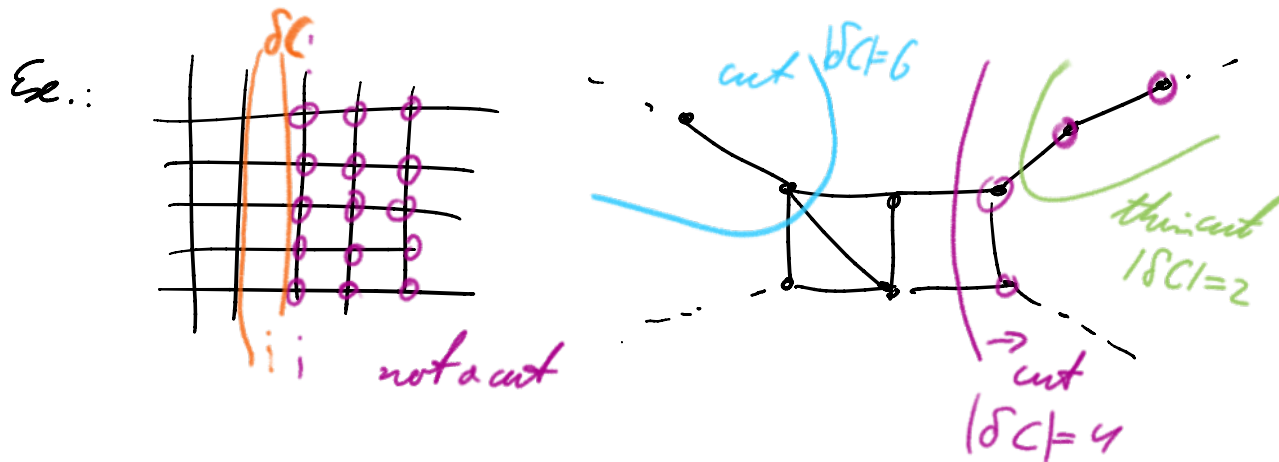
Hence, there is a path  $P$  from  $v$  to  $w$  in  $X - \{e\}$ . Every  $(k+1)$ -separator  $S$  in  $X$  that contains  $e$  also contains an edge  $e'$  of  $P$ . Furthermore,  $S \setminus \{e\}$  is a  $k$ -separator in  $X - \{e\}$ . By the induction hypothesis, there are only fin. many  $k$ -separators in  $X - \{e\}$  that <sup>contain</sup>  $e'$ . The statement follows because  $P$  is finite and different  $(k+1)$ -separators in  $X$  that <sup>contain</sup>  $e$  and  $e'$  correspond to different  $k$ -separators in  $X - e$  that contain  $e'$ .  $\square$

Def 10.7:

- A cut is a set of vertices  $C \subseteq X^0$  such that  $\delta C$  is finite and  $C$  and  $C^c$  both contain a ray.

• If there is a cut, define  $\chi := \min_{\text{cut}} \{|\delta C|\}$ .

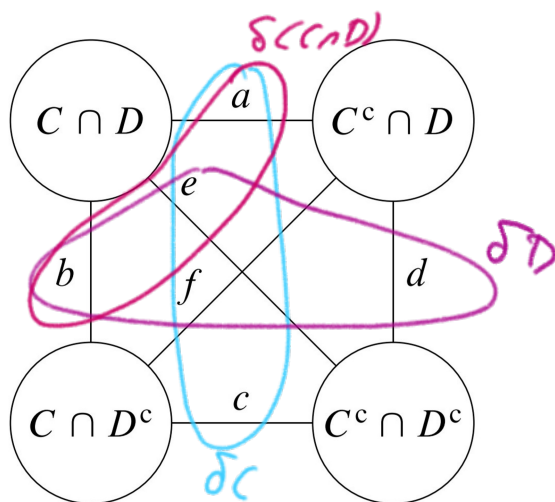
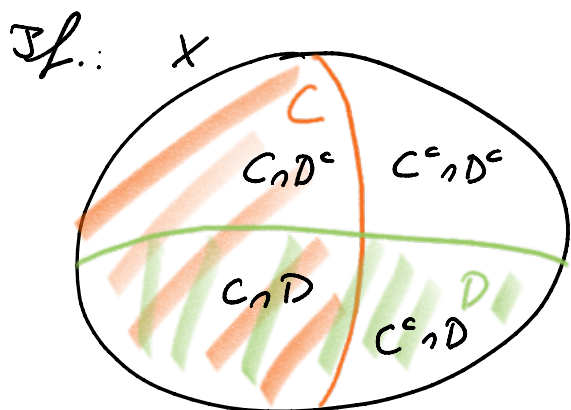
• A cut  $C$  with  $|\delta C| = \chi$  is called thin.



Rem.:

- Cuts are not necessarily connected. Thin cuts are.
- In connected graphs with more than one end, there is always a thin cut.

Lemma 10.8: Let  $C$  and  $D$  be thin cuts. If  $C \cap D$  and  $C^c \cap D^c$  are cuts, then they are thin.



Define

$$a = |\delta(C \cap D, C^c \cap D)|,$$

$$b = |\delta(C \cap D, C \cap D^c)|,$$

$$c = |\delta(C \cap D^c, C^c \cap D^c)|,$$

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$$d = |\delta(C^c \cap D, C^c \cap D^c)|,$$

$$e = |\delta(C \cap D, C^c \cap D^c)|,$$

$$f = |\delta(C \cap D^c, C^c \cap D)|.$$

Then

$$\begin{aligned} \chi &= |\delta C| = a + e + f + c \\ &= |\delta D| = b + e + f + d. \end{aligned}$$

Hence

$$2\chi = a + b + c + d + 2e + 2f. \quad (*_1)$$

Now if  $C \cap D$  and  $C^c \cap D^c$  are cuts, then

$$|\delta(C \cap D)| = a + b + c \geq \chi \text{ and}$$

$$|\delta(C^c \cap D^c)| = c + e + d \geq \chi.$$

$\Rightarrow a + b + c + d + 2e \geq 2\chi$  and <sup>so</sup> by  $(*_1)$ :

$$a + b + c + d + 2e = 2\chi \text{ and } f = 0.$$

Finally,  $a + e + b = c + e + d = \chi$

$$|\delta(\overline{C \cap D})| = |\delta(C^c \cap D^c)|.$$

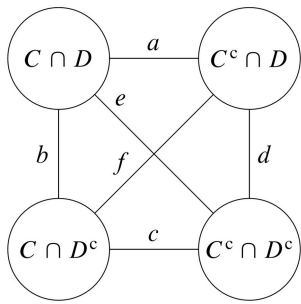
□

Def 10.9: Let  $C, D \subseteq X^0$ .

The intersections

$$C \cap D, C \cap D^c, C^c \cap D, C^c \cap D^c$$

are called the corners of  $C$  and  $D$ . We say that  $C \cap D$  is opposite to  $C^c \cap D^c$  and  $C \cap D^c$  is opposite to  $C^c \cap D$ .



- $C$  and  $D$  are nested if one of their corners is empty.
- $C$  is constant on  $D$  if either  $D \subseteq C$  or  $C \cap D = \emptyset$ .

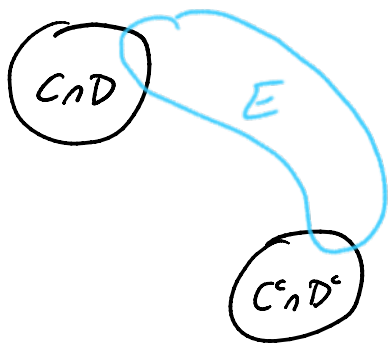
Note:  $C$  is nested with  $D$  iff  $C$  is constant on  $D$  or on  $D^c$ .

lem. 10.10: i) If there are two opposite corners of  $C$  and  $D$  such that  $E$  is not nested with either of them, then  $E$  is neither nested with  $C$  nor with  $D$ .

ii) Let  $C$  and  $D$  be cuts that are not nested. If a set  $E$  is nested with  $C$  and  $D$ , then  $E$  is nested with all corners of  $C$  and  $D$ .

3f:

i) If  $E$  is not nested with two opposite corners, then it is not constant on these corners. Hence it is not constant



on any of  $C, C^c, D, D^c$ .

( $E$  const. on  $C \Rightarrow E$  const. on any subset of  $C$ .)

ii) As  $C$  and  $D$  are not nested, all their corners are non-empty. After possibly replacing  $C$  with  $C^c$  and  $D$  with  $D^c$ , we can assume that  $E$  is constant on  $C$  and  $D$ . As  $C \cap D \neq \emptyset$ . This implies that  $E$  is const. on



w. can assume that  $E$  is constant on  $C$  and  $D$ .  
As  $C \cap D \neq \emptyset$ , this implies that  $E$  is const. on

$$C \cup D = (C \cap D) \cup (C^c \cap D) \cup (C \cap D^c).$$

Hence, for any corner  $A$ , the set is constant on  $A$  or  $A^c$ . □

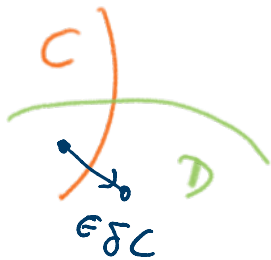
Def 10.11: Let  $C$  be a cut. We denote by  $M(C)$  the set of thin cuts that are not nested with  $C$ .

We set  $m(C) := |M(C)|$ .

Rem.: For all cuts  $C$ ,  $m(C)$  is finite.

I realised after the lecture that my explanations below were probably a bit short. For a more detailed proof that shows that for a thin cut  $C$ , the number  $m(C)$  is finite (this is all we need), see Dunwoody - "Cutting up graphs" 2.6.

$C$  not nested with  $D \Rightarrow D$  contains an edge of  $\delta C$



By Lemma 10.6, every edge of  $\delta C$  is only contained in fin many  $\delta D$ ,  $D$  is a thin cut.

Lemma 10.12: Let  $C$  and  $D$  be thin cuts that are not nested. If  $C \cap D$  and  $C^c \cap D^c$  are cuts, then

$$m(C \cap D) + m(C^c \cap D^c) < m(C) + m(D). \quad (*)$$

Prf.: By Lemma 10.8,  $C \cap D$  and  $C^c \cap D^c$  are thin.

Let  $E$  be a thin cut. If  $E$  is in  $M(C \cap D) \cap M(C^c \cap D^c)$ , then  $E$  is neither nested with  $C \cap D$  nor with  $C^c \cap D^c$ .

Hence by Lemma 10.10(i), it is in  $M(C)$  and  $M(D)$ .

This means that if  $E$  is counted twice on the LHS of  $(*)$ , then it is counted twice on the RHS of  $(*)$ .

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If  $E$  is counted once on the LHS of  $(*)$ , then it is in  $M(C \cap D) \setminus M(C^c \cap D^c)$  or in  $M(C^c \cap D^c) \setminus M(C \cap D)$ .

Then by Lem. 10.10 ii), it is in  $M(C)$  or  $M(D)$ , and it is counted at least once on the RHS of  $(*)$ .

This shows.

$$m(C \cap D) + m(C^c \cap D^c) \leq m(C) + m(D).$$

To see that the inequ. is strict, note that  $C \in M(D)$  but  $C \notin M(C \cap D)$  or  $M(C^c \cap D^c)$ , so it's counted on the RHS of  $(*)$ , but not on the LHS.  $\square$

Def 10.13: Let  $\mathcal{C}$  be the set of all thin cuts and set  $m := \min_{C \in \mathcal{C}} \{m(C)\}$ .

A thin cut  $C$  with  $m(C) = m$  is called optimally nested.

Rem.:  $\rightarrow$  As  $m(C)$  is finite for all  $C \in \mathcal{C}$ , the minimum is attained if  $\mathcal{C} \neq \emptyset$ . I.e. if there is a cut, there is an optimally nested cut.

Thm. 10.14: Optimally nested cuts are all nested with each other.

Pf.: Suppose  $C$  and  $D$  are optimally nested cuts that are not nested with each other. Then  $m \geq 1$



As  $C$  and  $D$  are cuts, there is a pair of opposite cones that contain a ray and thus are both cuts.

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If  $C \cap D$  and  $C^c \cap D^c$  are infinite, then they have an infinite connected component and hence contain a ray.

If e.g.  $C \cap D$  is finite, then  $C \cap D^c$  and  $C^c \cap D$  are infinite as  $C$  and  $D$  are inf.

After possibly replacing  $C$  with  $C^c$ , we may assume that  $C \cap D$  and  $C^c \cap D^c$  are cuts. But now Lemma 10.12 says that

$$m(C \cap D) + m(C^c \cap D^c) < m(C) + m(D) = 2m.$$

Hence, one of the summands on the LHS must be smaller than  $m$ , which is a contradiction.  $\square$