

# Machine Learning in Finance

## Exercise sheet 10

Through this exercise sheet, we let  $E = \mathbb{R}^d$ ,  $J$  an interval on  $\mathbb{R}$ , and denote  $\mathbf{Sig}_J: \mathcal{C}_0^1(J, E) \rightarrow \mathbf{T}(E)$  the signature map such that for all  $X \in \mathcal{C}_0^1(J, E)$ .

### Exercise 10.1 (Signatures and reservoirs computing)

(a) Let  $X \in \mathcal{C}_0^1([0, T], \mathbb{R}^n)$  satisfying the dynamic:

$$dX_t = \sum_{k=1}^m V_k(X_t) dW_t^k, \quad X_t \in \mathbb{R}^n, W_t \in \mathbb{R}^m, V^k: \mathbb{R}^n \rightarrow \mathbb{R}^n, \quad (1)$$

where  $(W_t)_{t=0}^\infty$  is a Brownian motion. Prove that

$$X_t = \sum_{d=0}^\infty \sum_{i_1, \dots, i_d=1}^n \left( \int_{0 \leq t_1 \leq \dots \leq t_d \leq t} dW_{t_1}^{i_1} \dots dW_{t_d}^{i_d} \right) V^{i_d} \dots V^{i_1}(X_0) \cdot X_0. \quad (2)$$

where

$$Vf(x) = df(x) \cdot V(x).$$

(b) Rewrite (2) with signature in the form of the following:

$$X_t = \langle \mathbf{R}, \mathbf{Sig}_{[0,t]}(W) \rangle X_0, \quad (3)$$

and express the readout  $\mathbf{R}$  with  $(V^k)_{k=1}^m$  (notice that  $\mathbf{R}$  depends on  $X_0$ ).

(c) Relate (3) with reservoirs computing.

### Exercise 10.2 (Randomized signature)

(a) Prove the Johnson-Lindenstrauss Lemma (see lecture notes): For any  $0 < \epsilon < 1$  and any integer  $n$ , let  $k$  be a positive integer such that

$$k \geq \frac{24}{3\epsilon^2 - 2\epsilon^3} \log n$$

then for any set  $A$  of at most  $n$  points in  $\mathbb{R}^d$ , there exists  $f: \mathbb{R}^d \rightarrow \mathbb{R}^k$  s.t. for all  $x, y \in A$  that

$$(1 - \epsilon)\|x - y\|^2 \leq \|f(x) - f(y)\|^2 \leq (1 + \epsilon)\|x - y\|^2. \quad (4)$$

Johnson-Lindenstrauss Lemma not only prove the existence but also provide the projection, what is the projection considered in the proof of Johnson-Lindenstrauss Lemma?

(b) Let  $X \in \mathcal{C}_0^1([0, T], E)$  and define  $\pi_n: \mathbf{T}(E) \rightarrow \mathbf{T}^{(n)}(E)$  the projection such that for all  $\mathbf{x} \in \mathbf{T}(E)$

$$\pi_n((\mathbf{x}_k)_{k \geq 0}) = (\mathbf{x}_k)_{k \leq n} \quad (5)$$

then  $S_t = \mathbf{Sig}_{[0,t]}^{(n)}(X)$  satisfies for all  $t \in [0, T]$  that

$$dS_t = \pi_n(S_t \otimes dX_t), \quad S_0 = (1, 0, \dots). \quad (6)$$

From this subproblem we can actually see signature as a solution of ODE which closely relates to the definition of randomized signature.

- (c) Recall that randomized signature is defined as the solution of the following dynamic:

$$dRS_t = \sum_{i=1}^d \sigma(A_i RS_t + b_i) dX_t^i, \quad RS_0 \in \mathbb{R}^k, \quad (7)$$

where  $A_i$  and  $b_i$  are pre-randomized and  $\sigma$  be a non-linear function. Let  $\sigma$  be real analytic with infinite radius of convergence and let  $A_1, \dots, A_d$  and  $b_1, \dots, b_d$  be independent samples of a probability law absolutely continuous with respect to Lebesgue measure. Then prove that linear combinations of randomized signature for all possible initial values  $RS_0 \in \mathbb{R}$  are dense in  $C(K)$  on compact subsets of bounded variation curves.

## References

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- [3] Christa Cuchiero, Lukas Gonon, Lyudmila Grigoryeva, Juan-Pablo Ortega, and Josef Teichmann. Discrete-time signatures and randomness in reservoir computing. *IEEE Transactions on Neural Networks and Learning Systems*, 2021.
- [4] Terry J Lyons, Michael Caruana, and Thierry Lévy. *Differential equations driven by rough paths*. Springer, 2007.