RANDOM WALKS ON TRANSITIVE GRAPHS (D-MATH) EXERCISE SHEET 9

Definition. Let \mathcal{X} be a finite set. Let \mathcal{B} the set of binary numbers. A binary code is a mapping $C : \mathcal{X} \to \mathcal{B}$. For $x \in \mathcal{X}$, let C(x) denote the codeword corresponding to x and let l(x) the length of C(x). We say that C is a binary prefix code if $\forall x, y \in \mathcal{X}$, C(x) is not a prefix of C(y) and C(y) is not a prefix of C(x).

For example, C(red) = 0, C(green) = 10, C(blue) = 11 is a binary prefix code for $\mathcal{X} = \{\text{red}, \text{green}, \text{blue}\}$ with codewords $\{0, 10, 11\}$. On the other hand C(red) = 0, C(green) = 0, C(blue) = 10 is a binary code but not a binary prefix code, since 0 is a prefix of 01.

(*) Exercise 1. (Kraft inequality) Let l_1, l_2, \ldots, l_m the lengths of the codewords of a binary prefix code $C : \mathcal{X} \to \mathcal{B}$, as defined above. We want to show that

$$\sum_{i} 2^{-l_i} \le 1. \tag{1}$$

(i) Consider a binary tree T of height $l_{\max} := \max\{l_1, \ldots, l_m\}$ where each vertex of the tree represents one of the codewords in the standard manner (root \emptyset , with children $\{0, 1\}$, each of those having children $\{00, 01\}$, and $\{10, 11\}$ respectively, and so on). Let U be a uniformly distributed random variable on the leaves of T. Let Z given by

 $\mathbb{P}[Z = C(x)] = \mathbb{P}[U \text{ contains a descendent of } C(x)], \text{ for all } x \in \mathcal{X}.$

Show that Z is a well defined random variable on $C(\mathcal{X})$ and calculate its law. Use this to show (1).

(ii) Show that the converse is also true, that is, given a set of codeword lengths that satisfy (1), there exists a binary prefix code with these word lengths.

(*) Exercise 2. Let X be a random variable with values in a finite set \mathcal{X} . Let $C : \mathcal{X} \to \mathcal{B}$ be a binary prefix code.

(i) Show that

$$\mathbb{E}[l(X)] \ge H_2(X) := -\sum_{x \in \mathcal{X}} \mathbb{P}[X = x] \log_2(\mathbb{P}[X = x]).$$

with equality if and only if $\mathbb{P}[X = x] = 2^{-l(x)}$, for all $x \in \mathcal{X}$. (ii) Show that $\mathbb{E}[l(X)] < H_2(X) + 1$. **Definition.** Let X be a random variable with values in a finite set \mathcal{X} . The Shannon entropy of X is defined by

$$H(X) := -\sum_{x \in \mathcal{X}} \mathbb{P}[X = x] \log(\mathbb{P}[X = x]).$$

If Y is a random variable with values in a finite set \mathcal{Y} , we define the joint entropy of (X, Y) by

$$H(X,Y) := -\sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} \mathbb{P}[X = x, Y = y] \log(\mathbb{P}[X = x, Y = y]).$$

Remark. We use the convention $0 \cdot \log(0) = 0$.

- (*) Exercise 3. Let X and Y as in the definition before. Show the following properties: (i) $H(X) \ge 0$.
 - (ii) $H(X) \leq \log(|\mathcal{X}|)$ with equality if and only if $X \sim \text{Uniform}(\mathcal{X})$. *Hint*: Use the fact that log is a concave function.
- (iii) $H(X,Y) \leq H(X) + H(Y)$ with equality if and only if X, Y are independent.
- (iv) $H(X) \leq H(X, Y)$ with equality if and only if Y is $\sigma(X)$ -measurable.

(*) Exercise 4. Let X and Y random variables in a common probability space $(\Omega, \mathcal{A}, \mathbb{P})$, taking values in finite sets \mathcal{X} and \mathcal{Y} , respectively. Let $\mathcal{F} \subset \mathcal{A}$ be a sub- σ -algebra. We define the *conditional entropy* of X given \mathcal{F} , by

$$H(X \mid \mathcal{F}) := -\mathbb{E}\left[\sum_{x \in \mathcal{X}} \mathbb{P}[X = x \mid \mathcal{F}] \log(\mathbb{P}[X = x \mid \mathcal{F}])\right].$$

- (i) Show that $H(X | \mathcal{F}) \leq H(X)$ with equality if and only if X is independent of \mathcal{F} .
- (ii) We define $H(X | Y) := H(X | \sigma(Y))$. Show that

$$H(X \mid Y) = -\sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} \mathbb{P}[X = x, Y = y] \log(\mathbb{P}[X = x \mid Y = y])$$

(iii) Show that $H(X \mid Y) = H(X, Y) - H(Y)$.